

Exploiting Branch Constraints without Exhaustive Path Enumeration

Ting Chen Tulika Mitra Abhik Roychoudhury Vivy Suhendra

School of Computing, National University of Singapore

{chent, tulika, abhik, vivy}@comp.nus.edu.sg

Abstract

Statically estimating the worst case execution time (WCET) of a program is important for real-time software. This is difficult even in the programming language level due to the inherent difficulty in detecting and exploiting infeasible paths in a program's control flow graph. In this paper, we propose an efficient method to exploit infeasible path information for WCET estimation of a loop without resorting to exhaustive path enumeration. The efficiency of our approach is demonstrated with a real-life control-intensive program.

1. Introduction

Static analysis of a program for obtaining the Worst Case Execution Time (WCET) is important for hard real-time embedded systems. WCET analysis typically consists of three phases: *flow analysis* to identify loop bounds and infeasible flows through the program; *architectural modeling* to determine the effect of pipeline, cache, branch prediction etc. on the execution time; and finally *estimation* to find an upper bound on the WCET of the program given the results of the flow analysis and the architectural modeling. In this paper, we concentrate on the estimation problem.

There exist mainly three different approaches for WCET estimation: *tree-based*, *path-based*, and *implicit path enumeration*. The tree-based approach estimates the WCET of a program through a bottom-up traversal of its syntax tree and applying different timing rules at the nodes (called "timing schema") [5]. This method is quite simple and efficient. But it has limitations in exploiting the results returned by flow analysis. In particular, it is difficult to exploit infeasible paths due to branch constraints (dependencies among branch statements) in this approach as the timing rules are local to a program statement. Implicit path enumeration techniques (IPET) [4] represent the program flows as linear equations or constraints and attempt to maximize the execution time of the entire program under these constraints. This is done via an Integer Linear Programming

(ILP) solver. Attempts have been made to integrate special flow information in IPET [2]. However, the kind of flow information that can be handled by IPET is inherently limited. This is because the usual ILP formulation introduces formal variables for the execution counts of the nodes and edges in the Control Flow Graph (CFG) of the program. Since the variables denote aggregate execution counts of basic blocks, it is not possible to express certain infeasible path patterns (typically denoting a *sequence* of basic blocks) as constraints on these variables.

Path-based techniques estimate the WCET by computing execution time for the feasible paths in the program and then searching for the one with the longest execution time. Thus, path-based techniques can naturally handle the various flow information. Healy et al., in particular, detect and exploit branch constraints within the framework of path-based technique [3]. Originally, path-based techniques were limited to a single loop iteration. However, Stappert et al. [6] have extended it to complex programs with the help of *scope graphs* and *virtual scopes*.

One of the main drawbacks of path-based techniques is that they require the generation of all the paths. In the worst case, this can lead to 2^n paths where n is the number of decisions in the program fragment. In control-intensive programs, we have encountered up to 6.55×10^{16} paths in a single loop iteration (see Table 1). Research by Stappert et al. [6] has sought to avoid this expensive path enumeration by finding (a) the longest program path, (b) checking for the feasibility of, and (c) removing from CFG followed by the search for a new longest path if is infeasible. This technique is a substantial improvement over exhaustive path enumeration. However, if the feasible paths in the program have relatively low execution times, then this approach still has to examine many program paths. Indeed, for our benchmark only a small fraction (less than 0.1%) of the paths are feasible, making this approach quite costly (see Table 1).

In this paper, we present a technique for finding the WCET of a program in the presence of infeasible paths without performing exhaustive path enumeration.

2. WCET Estimation Algorithm

In this section, we present a method for finding the WCET path of a single loop. Once this is obtained, the WCET path of the program can be obtained by composing the WCET paths of individual loops through the well-known timing schema approach [5]. Of course, this will mean that infeasible path information across loops cannot be taken into account. We assume that state-of-the-art WCET analyzers such as aiT[1] can be used to estimate the worst-case execution time of each basic block.

Given a fragment of assembly code corresponding to a loop in a source program, we first construct the directed acyclic graph (DAG) capturing the control flow in the loop body (i.e., the CFG of the loop body without the loop back-edge). We assume that the DAG has a unique source node and a unique sink node. If there is no unique sink node, then we add a dummy sink node. Each path from the source to the sink in the DAG is an **acyclic path** — a possible path in a loop iteration. Our algorithm finds the worst case execution path for a single iteration of the loop, i.e., we find the heaviest acyclic path. If the estimated execution time of the heaviest acyclic path is t and the loop-bound is lb , then the loop's estimated WCET is $lb \cdot t$.

We find the heaviest acyclic path accurately by taking into account infeasible path information and yet we avoid enumerating all the acyclic paths in a loop. Clearly, the infeasible path information that we work with may not always be complete; so the accuracy of our heaviest acyclic path detection depends on the accuracy of the infeasible path information. First, we discuss the infeasible path information used and then explain how it is efficiently exploited in our WCET calculation.

2.1. Infeasible path information

Our infeasible path information consists of two binary relations capturing conflicting pairs of branches/assignments: *AB_conflict* and *BB_conflict*. The relation *AB_Conflict* is a set of (assignment, branch-edge) pairs, that is, if $a, e \in AB_Conflict$ then a is an assignment instruction and e is an outgoing edge from a conditional branch instruction; on the other hand, the relation *BB_Conflict* is a set of (branch-edge, branch-edge) pairs.

We do not detect conflicts between arbitrary branches and assignments to avoid an inefficient conflict detection procedure. The only conditional branches whose edges appear in our *BB_conflict* and *AB_conflict* relations are of the form *variable relational_operator constant*. Similarly, the only assignments which appear in *AB_Conflict* are of the form *variable := constant*. For such assignments and branches we can define pair-wise conflict in a natural way

(see [3] for a full discussion). For example, $x := 2$ conflicts with $x > 3$, but not with $x < 3$; similarly $x > 3$ conflicts with $x < 2$ but not with $x > 5$. Now, for such restricted branches and assignments, we put an assignment a conflicting with a conditional branch-edge e into the *AB_Conflict* relation (i.e., $a, e \in AB_Conflict$) iff there exists at least one path from a to e which does not contain assignments to the common variable appearing in a, e . Similarly, we put two conflicting conditional branch-edges e_1, e_2 into the *BB_Conflict* relation (i.e., $e_1, e_2 \in BB_Conflict$) iff there exists at least one path from e_1 to e_2 which does not contain assignments to the common variable appearing in e_1, e_2 . If $e_1, e_2 \in BB_Conflict$, there may be another path from e_1 to e_2 that contains an assignment to the common variable appearing in e_1, e_2 . Such paths should not be considered as infeasible paths.

The computation of the *AB_Conflict* and *BB_Conflict* relations can be accomplished in $O((|V| + |E|) \cdot |E|)$ time where $|V|, |E|$ are the number of nodes and the number of edges in the control flow DAG; this is because for each branch-edge we need to perform a depth-first like search to find conflicting branch-edges and/or assignments.

Example: A loop-free program fragment (which can be the body of a loop) and its control flow DAG are shown in Figure 1. In this example, the relation *AB_Conflict* contains only one pair – the assignment at basic block $B6$ (which sets x to 1) and the branch-edge $B7 \rightarrow B9$ (which stands for $x > 2$). The relation *BB_Conflict* contains the branch-edge pair $B1 \rightarrow B2, B7 \rightarrow B8$ that captures the conditions $x > 3$ and $x < 2$.

2.2. WCET calculation

We now present our WCET estimation algorithm for finding the heaviest feasible path in an iteration of a loop. We *do not* enumerate the possible paths in an iteration and then find the heaviest. At the same time, we do not consider all paths in the loop's control flow graph to be feasible – we consider the infeasible path information captured by *AB_Conflict* and *BB_Conflict* relations.

Our algorithm traverses the loop's control flow DAG from sink to source. However, to take into account the infeasible path information, we cannot afford to remember only the "heaviest path so far" as we traverse the DAG. This is because the heaviest path may have conflicts with earlier branch-edges or assignment instructions resulting in costly backtracking. Instead, at a basic block v , we maintain a set of paths $paths(v)$ where each $p \in paths(v)$ is a path from v to the sink node. $paths(v)$ contains only those paths which when extended up to the source node can potentially become the WCET path. For each path $p \in paths(v)$ we also maintain a "conflict list". The conflict list contains

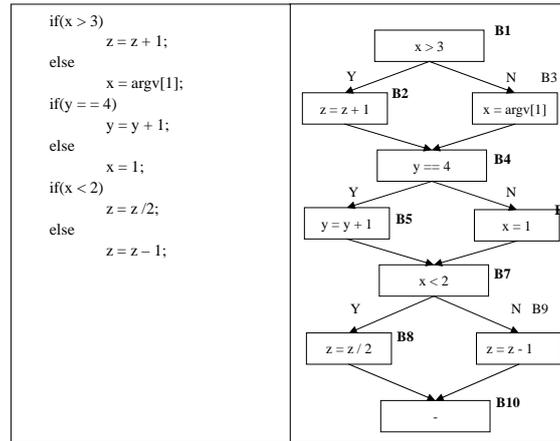


Figure 1: A loop body with its Control Flow Graph

the branch-edges of p that participate in conflict with ancestor nodes and edges of v .

During our traversal from sink to source, consider a single step traversal from v to u along the edge $u \rightarrow v$ in the control flow DAG. We first construct $paths(u)$ from $paths(v)$ by adding the edge $u \rightarrow v$ at the beginning of the paths in $paths(v)$. Also, for each path $p \in paths(u)$ we update its conflict list to contain exactly those edges in p which have conflicts with branch-edges/assignments “prior to” u (in topological order). At this stage we may add the branch-edge $u \rightarrow v$ to p ’s conflict list or we may remove an edge e from p ’s conflict list if all branch-edges/assignments conflicting with e appear “after” u (in the topological order). If as a result, we have identical conflict lists for two paths $p, p' \in paths(u)$, then we maintain the heavier path among p and p' . Finally, if the conflict list of a path $p \in paths(u)$ becomes empty and p is the heaviest path in $paths(u)$, we assign the singleton set $\{p\}$ to $paths(u)$. Details of the algorithm are omitted for space considerations.

In the worst case, the complexity of our algorithm is exponential in $|V|$, the number of nodes in the loop’s control flow DAG. This is because the number of paths in $paths(v)$ for some block v may be $O(2^{|V|})$ due to different decisions in the branches following v . In practice, this exponential blow-up is not encountered because (a) branch-edges which do not conflict with any assignment/branch-edge do not need to be kept track of, and (b) a branch-edge which conflicts with other branch-edges/assignments need not be remembered after we encounter those conflicting branch-edges/assignments during the traversal.

Illustration We demonstrate our WCET calculation method by employing it on the control flow DAG of Figure 1. As mentioned, we traverse the DAG from sink to source and maintain a set of paths $paths(v)$ at each visited node v . For each path $p \in paths(v)$ we also maintain the *conflict list*, a set of branch decisions drawn from branch decisions made

so far. Thus each path p in $paths(v)$ is written in the form $p \text{ conflict list}$.

Starting from node B_{10} in Figure 1, our traversal is routine till we reach node B_7 (\emptyset denotes empty set).

$$\begin{aligned}
 paths(B_{10}) &= \{ B_{10} \} \\
 paths(B_9) &= \{ B_9, B_{10} \} \\
 paths(B_8) &= \{ B_8, B_{10} \}
 \end{aligned}$$

The outgoing edges from node B_7 appear in conflict relations capturing infeasible path information. Consequently, our method maintains two paths in $paths(B_7)$ — the heaviest path starting with $B_7 \rightarrow B_8$ and the heaviest path starting with $B_7 \rightarrow B_9$.

$$\begin{aligned}
 paths(B_7) &= \{ B_7, B_8, B_{10} \}_{B_7 \rightarrow B_8} \\
 &\quad B_7, B_9, B_{10} \}_{B_7 \rightarrow B_9}
 \end{aligned}$$

Now, from node B_7 we traverse to nodes B_5 and B_6 . The assignment in node B_6 conflicts with $B_7 \rightarrow B_9$. Therefore, we do not consider any path in $paths(B_7)$ which contains $B_7 \rightarrow B_9$ in its conflict list. This is how infeasible path information is accounted for in our WCET calculation. Thus we have

$$paths(B_6) = \{ B_6, B_7, B_8, B_{10} \}$$

We drop $B_7 \rightarrow B_8$ from the conflict list of B_6, B_7, B_8, B_{10} as we have encountered an assignment to program variable x in B_6 . The assignment implies that the conflict between $B_7 \rightarrow B_8$ and $B_1 \rightarrow B_2$ does not hold along any extension of the partial path B_6, B_7, B_8, B_{10} .

At node B_5 , we first add B_5 to the two partial paths from B_7 . Then, we notice that the edge $B_7 \rightarrow B_9$ is involved only in a conflict with B_6 and we have already traversed B_6 . Therefore, we can drop this edge from the conflict list of the partial path B_5, B_7, B_9, B_{10} and this path now becomes completely conflict free. Assuming that

| Function | Basic Blocks | Total Paths | Feasible paths | BB-Conflicts | AB-Conflicts | Enumerated Paths |
|------------|--------------|-----------------------|-----------------------|--------------|--------------|------------------|
| statemate | 334 | 6.55×10^{16} | 1.09×10^{13} | 74 | 15 | 121,831 |
| statemate1 | 44 | 19,440 | 7,440 | 6 | 0 | 15 |
| statemate2 | 73 | 902 | 36 | 26 | 0 | 14 |
| statemate3 | 161 | 1,459,364 | 69,867 | 15 | 0 | 40 |
| statemate4 | 17 | 10 | 10 | 0 | 0 | 1 |
| statemate5 | 43 | 256 | 58 | 2 | 0 | 4 |

Table 1: Efficiency of our WCET calculation method.

| Function | WCET Estimation (cycles) | |
|------------|--------------------------|--------------------|
| | w/o infeasibility | with infeasibility |
| statemate | 44,800 | 41,520 |
| statemate1 | 29,400 | 28,960 |
| statemate2 | 2,750 | 2,270 |
| statemate3 | 7,300 | 7,000 |
| statemate4 | 1,070 | 1,070 |
| statemate5 | 2,370 | 2,090 |

Table 2: Accuracy of WCET estimation with and without considering infeasibility.

B_5, B_7, B_9, B_{10} is heavier than B_5, B_7, B_8, B_{10} , we have

$$paths(B_5) = \{ B_5, B_7, B_9, B_{10} \}$$

On the other hand, if B_5, B_7, B_8, B_{10} is heavier, we still need to maintain both the partial paths (as the heavier path may become infeasible later). Continuing in this way we reach node B_1 ; we omit the details for the rest of the traversal. Note that the control flow DAG of Figure 1 has three branches and $2^3 = 8$ paths. However, when we visit any basic block v of the control flow DAG, $paths(v)$ contains at most two paths (*i.e.*, the exponential blow-up is avoided here in practice).

3. Experiments

In this section, we present preliminary experiment results for the proposed method. The benchmark used in our experiment is a car window lift control program taken from the C-Lab benchmark suite. It is automatically generated by the Statemate tool based on a state-chart specification. We report results for the entire program (statemate) as well as program fragments corresponding to all the five functions (statemate1 to statemate5).

An enumeration-based WCET estimation method typically examines each possible path, filters out the infeasible paths and selects the feasible path with the maximum execution time. Table 1 shows that the total number of paths through a single iteration of the loop body can be quite large (6.55×10^{16} possible paths out of which at most 1.09×10^{13} are feasible for statemate). In fact, a naive WCET calculation method, which enumerates all these possible paths

for one loop iteration and chooses the longest feasible one, runs out of memory even on a PC with 1 GB main memory. The column *Enumerated Paths* shows the maximum number of paths that need to be maintained by our estimation technique at any point of time. The results are quite encouraging; even for the entire statemate program, we only need to keep at most 121,831 paths at any point of time during the estimation. As a result, our estimating technique requires less than 1 minute for the entire statemate program on a Pentium4 1.7Ghz platform with 1GB memory. Finally, Table 2 shows that, as expected, our method produces more accurate estimation compared to a method that does not take infeasibility information into account. The only exception is statemate4, which does not have any infeasible path.

4. Discussion

In this paper, we have reported preliminary results on exploiting (limited) infeasible path information during WCET estimation of a loop without resorting to path enumeration. The efficacy of our technique has been demonstrated on a substantial real-life car window control benchmark. In near future, we will develop WCET estimation methods that can take into account infeasible path patterns of arbitrary length without compromising efficiency.

References

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