Solving Homogeneous Linear Equations over Polynomial Semirings

Ruiwen Dong \square

Department of Computer Science, University of Oxford, UK

— Abstract

For a subset B of \mathbb{R} , denote by U(B) be the semiring of (univariate) polynomials in $\mathbb{R}[X]$ that are strictly positive on B. Let $\mathbb{N}[X]$ be the semiring of (univariate) polynomials with non-negative integer coefficients. We study solutions of homogeneous linear equations over the polynomial semirings U(B) and $\mathbb{N}[X]$. In particular, we prove local-global principles for solving single homogeneous linear equations over these semirings. We then show PTIME decidability of determining the existence of non-zero solutions over $\mathbb{N}[X]$ of single homogeneous linear equations.

Our study of these polynomial semirings is largely motivated by several semigroup algorithmic problems in the wreath product $\mathbb{Z} \wr \mathbb{Z}$. As an application of our results, we show that the Identity Problem (whether a given semigroup contains the neutral element?) and the Group Problem (whether a given semigroup is a group?) for finitely generated sub-semigroups of the wreath product $\mathbb{Z} \wr \mathbb{Z}$ is decidable when elements of the semigroup generator have the form $(y, \pm 1)$.

2012 ACM Subject Classification Computing methodologies \rightarrow Symbolic and algebraic manipulation

Keywords and phrases wreath product, identity problem, polynomial semiring, positive polynomial

Digital Object Identifier 10.4230/LIPIcs.STACS.2023.26

Funding The author acknowledges support from UKRI Frontier Research Grant EP/X033813/1.

Acknowledgements The author would like to thank Markus Schweighofer for useful discussions and feedback and for pointing out the references [24] and [25].

1 Introduction

Linear equations over semirings appear in various domains in mathematics and computer science, such as automata theory, optimization, and algebra of formal processes [2, 3, 6, 11]. There have been numerous studies on linear equations over different semirings [12], for example the semiring of natural numbers (integer programming), tropical semirings [23] and polynomial semirings [9, 22]. Given a semiring S, define S[X] to be the set of polynomials in variable X whose coefficients are elements of S. The set S[X] is again a semiring. One of the simplest polynomial semirings is the semiring $\mathbb{N}[X]$ of single variable polynomials with non-negative integer coefficients. The problem of solving a system of linear equations over $\mathbb{N}[X]$ was shown to be undecidable by Narendran [22] using a reduction from Hilbert's tenth problem. More precisely, given integer polynomials $h_{ij}, g_j \in \mathbb{Z}[X], i = 1, \ldots, n, j = 1, \ldots, k$, it is undecidable whether the system of equations

$$f_1h_{1j} + \dots + f_nh_{nj} = g_j, \quad j = 1, \dots, k,$$
(1)

has a solution (f_1, \ldots, f_n) over $\mathbb{N}[X]$. This contrasts with the decidability of solving systems of linear equations over \mathbb{N} and over $\mathbb{Z}[X]$ (using respectively integer programming [15] and Smith canonical forms [16]).

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Leibniz International Proceedings in Informatics

LIPICS Schloss Dagstuhl – Leibniz-Zentrum für Informatik, Dagstuhl Publishing, Germany

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In this paper, we show a decidability result for finding a non-zero solution of a *single* homogeneous linear equation over $\mathbb{N}[X]$. In particular, we are concerned with the following problem: given integer polynomials $h_1, \ldots, h_n \in \mathbb{Z}[X]$, does the equation

$$f_1h_1 + \dots + f_nh_n = 0 \tag{2}$$

admit a solution (f_1, \ldots, f_n) over $\mathbb{N}[X] \setminus \{0\}$ (i.e. none of the f_i is zero)?

In Section 6 of this paper we give a PTIME algorithm that decides this problem. Our algorithm relies on a local-global principle which we prove in Section 5, and reduces the decision problem to the existential theory of the reals in one variable. Formal definitions of these results will be given in Section 2.

It turns out that the problem of solving linear equations over the semiring $\mathbb{N}[X]$ is closely related to solving the same equation over the semiring U(B), consisting of polynomials in $\mathbb{R}[X]$ that are strictly positive on a subset B of \mathbb{R} . It is also related to the semiring W(B)of polynomials that are *non-negative* on B. The characterization of polynomials in U(B)and W(B) is a central subject in the theory of real algebra. In particular, when B is a semialgebraic set, variants of the *positivstellensatz* give explicit descriptions of the semirings U(B) and W(B). This theory can be traced back to the celebrated Hilbert's seventeenth problem: given a polynomial that takes only non-negative values over the reals, can it be represented as a sum of squares of rational functions? This has been answered positively by Artin [1] using a model theoretic approach. The techniques proposed by Artin have since developed into the rich theory of real algebra; for a comprehensive account of this subject, see [24] or [25]. An important result in solving homogeneous linear equations over $W(\mathbb{R})$ is the Bröcker-Prestel's local-global principle for weak isotropy of quadratic forms [24, Theorem 8.12, 8.13]. Applied over the function field $\mathbb{R}(X)$, the Bröcker-Prestel local-global principle relates the existence of non-trivial solutions over sums of squares in $\mathbb{R}(X)$ (and hence over $W(\mathbb{R})$ of a homogeneous linear equation, to the behaviour of the equation in all Henselizations of $\mathbb{R}(X)$. In Section 4 of this paper we prove a "strictly positive" version of the Bröcker-Prestel local-global principle, which characterizes the existence of solutions over U(B). This will serve as a base for proving further results in Section 5 and 6. Our proof is inspired by Prestel's proof of the original theorem. However, several new ideas are introduced to deal with the strict positivity as well as the positivity constraint over a subset of \mathbb{R} .

An important motivation for studying linear equations over $\mathbb{N}[X]$ comes from a semigroup algorithmic problem in the *wreath product* $\mathbb{Z} \wr \mathbb{Z}$. The wreath product is a fundamental construction in group and semigroup theory. Given two groups G and H, their wreath product $G \wr H$ is defined in the following way. Let G^H be the set of all functions $y: H \to G$ with finite support; it is a group with respect to pointwise multiplication. The group H acts on G^H as a group of automorphisms: if $h \in H, y \in G^H$, then $y^h(b) = y(bh^{-1})$ for all $b \in H$. The wreath product $G \wr H$ is then defined as the semi-direct product $G^H \rtimes H$, that is, the set of all pairs (y,h) where $y \in G^H, h \in H$, with multiplication operation given by

$$(y,h)(z,k) = (y^k z, hk).$$

One easy way to understand the group $\mathbb{Z} \wr \mathbb{Z}$ is through its isomorphism to a matrix group over the Laurent polynomial ring $\mathbb{Z}[X, X^{-1}]$ [19]:

$$\varphi \colon \mathbb{Z} \wr \mathbb{Z} \xrightarrow{\sim} \left\{ \begin{pmatrix} 1 & f \\ 0 & X^b \end{pmatrix} \middle| f \in \mathbb{Z}[X, X^{-1}], b \in \mathbb{Z} \right\}, \quad (y, b) \mapsto \begin{pmatrix} 1 & \sum_{k \in \mathbb{Z}} y(k) X^k \\ 0 & X^b \end{pmatrix}.$$
(3)

A large number of important groups are constructed using the wreath product, such as the lamplighter group $\mathbb{Z}_2 \wr \mathbb{Z}$ [13] and groups resulting from the Magnus embedding theorem [19]. The wreath product also plays an important role in the algebraic theory of automata. The Krohn–Rhodes theorem states that every finite semigroup (and correspondingly, every finite automaton) can be decomposed into elementary components using wreath products [17].

In Section 7 we give an application of our results to the *Identity Problem* in $\mathbb{Z} \wr \mathbb{Z}$. Given a finite set of elements $\mathcal{G} = \{A_1, \ldots, A_k\}$ in a group G as well as a target element $A \in G$, denote by $\langle \mathcal{G} \rangle$ the semigroup of generated by \mathcal{G} , and by $\langle \mathcal{G} \rangle_{grp}$ the group generated by \mathcal{G} . Consider the following decision problems:

- (i) (Group Membership Problem) whether $A \in \langle \mathcal{G} \rangle_{grp}$?
- (ii) (Semigroup Membership Problem) whether $A \in \langle \mathcal{G} \rangle$?
- (iii) (Identity Problem) whether the neutral element I of G is contained in $\langle \mathcal{G} \rangle$?

All three problems remain undecidable even when the ambient group G is restricted to relatively simple groups, such as the direct product $F_2 \times F_2$ of two free groups over two generators [5, 21]. Indeed, one of the first undecidability results in algorithmic theory was the undecidability of the Semigroup Membership Problem for integer matrices, obtained by Markov [20]. Some decidability results for the Identity Problem include its NP-completeness in $SL(2,\mathbb{Z})$ [4] and its PTIME decidability in nilpotent groups of class at most ten [10].

Let $p \in \mathbb{Z}_{>0}$. The group $\mathbb{Z} \wr \mathbb{Z}$ shares some common properties with the wreath product $(\mathbb{Z}/p\mathbb{Z}) \wr \mathbb{Z}$ and with the Baumslag-Solitar group $\mathsf{BS}(1,p)$. Similar to the isomorphism (3), both $(\mathbb{Z}/p\mathbb{Z}) \wr \mathbb{Z}$ and $\mathsf{BS}(1,p)$ can be represented as 2×2 upper triangular matrix groups:

$$\begin{aligned} (\mathbb{Z}/p\mathbb{Z}) \wr \mathbb{Z} &\cong \left\{ \begin{pmatrix} 1 & f \\ 0 & X^b \end{pmatrix} \middle| f \in (\mathbb{Z}/p\mathbb{Z}) [X, X^{-1}], b \in \mathbb{Z} \right\}, \\ \mathsf{BS}(1, p) &\cong \left\{ \begin{pmatrix} 1 & f \\ 0 & p^b \end{pmatrix} \middle| f \in \mathbb{Z}[1/p], b \in \mathbb{Z} \right\}. \end{aligned}$$

Lohrey, Steinberg and Zetzsche showed decidability of the Rational Subset Membership Problem (which subsumes all three decision problems mentioned above) in $H \wr V$, where His a finite and V is virtually free [18]. This notably implies its decidability in $(\mathbb{Z}/p\mathbb{Z}) \wr \mathbb{Z}$. Cadilhac, Chistikov and Zetzsche proved its decidability in BS(1, p) [7]. For $\mathbb{Z} \wr \mathbb{Z}$, decision problems are much harder due to higher encoding power of the ring $\mathbb{Z}[X, X^{-1}]$. The Group Membership Problem in $\mathbb{Z} \wr \mathbb{Z}$ can be reduced to the membership problem for modules over the ring $\mathbb{Z}[X, X^{-1}]$, and is hence decidable [26]. As for the Semigroup Membership Problem in $\mathbb{Z} \wr \mathbb{Z}$, Lohrey *et al.* showed its undecidability using an encoding of 2-counter machines [18]. Decidability of the Identity Problem in $\mathbb{Z} \wr \mathbb{Z}$ remains an intricate open problem. In this paper we give a decidability result in the case where all the elements of the generator \mathcal{G} are of the form $(y, \pm 1)$.

2 Main results

In this section we sum up the main results of this paper. For a subset B of \mathbb{R} , denote by U(B) the set of polynomials in $\mathbb{R}[X]$ that are strictly positive on B:

$$U(B) \coloneqq \{ f \in \mathbb{R}[X] \mid f(x) > 0 \text{ for all } x \in B \}.$$

Define \overline{B} to be the closure of B in \mathbb{R} under the Euclidean topology. Our first result is a local-global principle for solutions of homogeneous linear equations over U(B). Theorem 2.1 will be proved in Section 4.

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▶ **Theorem 2.1.** Let *B* be a subset of \mathbb{R} . Assume that $h_1, \ldots, h_n \in \mathbb{R}[X]$ are polynomials that satisfy $gcd(h_1, \ldots, h_n) = 1$. If the equation $f_1h_1 + \cdots + f_nh_n = 0$ has no solution (f_1, \ldots, f_n) over U(B), then there exists a real number $t \in \overline{B}$, such that the values $h_i(t), i = 1, \ldots, n$, are either all non-negative or all non-positive.

Our second result is a corollary of the previous theorem, it provides a similar local-global principle for solutions over $\mathbb{N}[X] \setminus \{0\}$. Theorem 2.2 will be proved in Section 5.

▶ **Theorem 2.2.** Given polynomials $h_1, \ldots, h_n \in \mathbb{Z}[X]$ with $gcd(h_1, \ldots, h_n) = 1$. If the equation $f_1h_1 + \cdots + f_nh_n = 0$ has no solution (f_1, \ldots, f_n) over $\mathbb{N}[X] \setminus \{0\}$, then there exists $t \in \mathbb{R}_{\geq 0}$, such that the values $h_i(t), i = 1, \ldots, n$ are either all non-negative or all non-positive.

Our next result shows that it is decidable in PTIME whether a linear homogeneous equation is solvable over $\mathbb{N}[X] \setminus \{0\}$. The input size is defined as the total number of bits used to encode all the coefficients of all h_i . Theorem 2.3 will be proved in Section 6.

▶ **Theorem 2.3.** Given as input $h_1, \ldots, h_n \in \mathbb{Z}[X]$. It is decidable in polynomial time whether the equation $f_1h_1 + \cdots + f_nh_n = 0$ has solutions f_1, \ldots, f_n over $\mathbb{N}[X] \setminus \{0\}$.

An application of this theorem is the following partial decidability result on the Identity Problem in the wreath product $\mathbb{Z} \wr \mathbb{Z}$. This will be the main topic of Section 7.

▶ **Theorem 2.4.** Given a finite set of elements $\mathcal{G} = \{(y_1, b_1), \dots, (y_n, b_n)\}$ in $\mathbb{Z} \wr \mathbb{Z}$, where $b_i = \pm 1$ for all *i*. The following are decidable:

1. (Group Problem) whether the semigroup $\langle \mathcal{G} \rangle$ generated by \mathcal{G} is a group.

2. (Identity Problem) whether the neutral element I is in the semigroup $\langle \mathcal{G} \rangle$.

3 Preliminaries

In this section we introduce the necessary mathematical tools on (semi)orderings of fields as well as valuations. Most notations and definitions follow those given in Prestel's book [24].

3.1 Orderings and semiorderings

▶ Definition 3.1 (Ordering). A linear ordering of a set S is a binary relation that satisfies
(i) a ≤ a,

- (ii) $a \leq b, b \leq c \implies a \leq c$,
- (iii) $a \leq b, b \leq a \implies a = a$,
- (iv) $a \leq b$ or $b \leq a$
- for all $a, b, c \in S$.

Given a field F of characteristic zero, a *(field) ordering* of F is a linear ordering \leq of the underlying set of F that additionally satisfies

(i) $a \le b \implies a + c \le b + c$,

(ii)
$$0 \le a, 0 \le b \implies 0 \le ab$$

for all $a, b, c \in F$. A field is called *formally real* if it admits at least one ordering.

The semiordering of a field, defined below, is a weaker version of the field ordering.

▶ Definition 3.2 (Semiordering). A semiordering of a field F is a linear ordering \leq of the underlying set of F that satisfies

(i) $a \le b \implies a + c \le b + c$, (ii) $0 \le 1$, (iii) $0 \le a \implies 0 \le ab^2$ for all $a, b, c \in F$. In a field F with semiordering \leq , we have $0 \leq x^2$ for all $x \in F$. The field of real numbers \mathbb{R} hence admits a unique semiordering, since every positive real can be written as a square. This semiordering is simply the natural ordering on \mathbb{R} .

It is easy to see that an ordering is always a semiordering. Conversely, a semiordering need not be an ordering. However, in any field, the existence of a semiordering implies that of an ordering.

Lemma 3.3 ([24, Corollary 1.15]). A field F is formally real (admits an ordering) if and only if it admits a semiordering.

For any subset P of F, define $-P := \{-x \mid x \in P\}$. For a semiordering \leq of F, the set $P := \{a \in F \mid 0 \leq a\}$ satisfies

- (i) $P + P \subseteq P$,
- (ii) $F^2 \cdot P \subseteq P$ and $1 \in P$,
- (iii) $P \cap -P = \{0\},\$
- (iv) $P \cup -P = F$.

Such a set will be called a *semicone* of F. A semicone P of F determines a semiordering \leq of F by $a \leq b \iff b - a \in P$. Therefore, we will sometimes call P a semiordering as well.

The *pre-semicone* is yet a weaker version of the semiordering (or semicone).

▶ Definition 3.4 (Pre-semicone). A *pre-semicone* of a field F is a subset P of F that satisfies (i) $P + P \subseteq P$,

(ii) $F^2 \cdot P \subseteq P$,

(iii) $P \cap -P = \{0\}.$

The only difference between a pre-semicone and a semicone is the absence of the rule (iv) and the condition $1 \in P$ in (ii). Obviously every semicone is also a pre-semicone. Conversely, a pre-semicone need not be a semicone, but it can always be extended to one.

▶ Lemma 3.5 ([24, Lemma 1.13]). For every pre-semicone P_0 of a formally real field F there exists a set $P \supseteq P_0$ such that P or -P is a semicone of F.

Suppose F is of characteristic zero. A semiordering or an ordering \leq of F is called *archemedean* if for each $a \in F$ there exists $n \in \mathbb{N} \subseteq F$ such that $a \leq n$.

▶ Lemma 3.6 ([24, Lemma 1.20]). Every archimedean semiordering is an ordering.

3.2 Valuations

Let F be a field. A valuation of F is a surjective map $v: F \to \Gamma \cup \{\infty\}$, where the value group Γ is an abelian totally ordered group¹, such that the following conditions are satisfied for all $a, b \in F$:

(i) $v(a) = \infty$ if and only if a = 0,

(ii) v(ab) = v(a) + v(b),

(iii) $v(a+b) \ge \min\{v(a), v(b)\}$, with equality if $v(a) \ne v(b)$.

A valuation is called *non-trivial* if $\Gamma \neq \{0\}$. A valued field is a pair (F, v) where F is a field and v is a valuation of F. Its valuation ring A_v is defined as

 $A_v \coloneqq \{a \in F \mid v(a) \ge 0\}.$

¹ An abelian totally ordered group Γ is an abelian group equipped with a linear ordering \leq , such that $a \leq b \implies a + c \leq b + c$ for all $a, b, c \in \Gamma$. Here, the group law of Γ is written additively. The ordering and the group law on Γ can be extended to the set $\Gamma \cup \{\infty\}$ by defining $a \leq \infty$ and $a + \infty = \infty + a = \infty + \infty = \infty$ for all $a \in \Gamma$.

We have $A_v \neq F$ if and only if v is non-trivial. A_v is a ring with a unique maximal ideal

 $M_v \coloneqq \{a \in F \mid v(a) > 0\}.$

The quotient $F_v := A_v/M_v$ is called the *residue field* of (F, v). It is indeed a field since M_v is maximal. A valuation v is called a *real place* of F if the residue field F_v is formally real.

Consider the field $F = \mathbb{R}(X)$. The following proposition gives a well-known characterization (up to isomorphism of the value group Γ) of the set of all non-trivial real places $\mathbb{R}(X)$ whose valuation ring contains the subfield \mathbb{R} .

▶ **Proposition 3.7.** Let v be a non-trivial real place of $\mathbb{R}(X)$ such that $\mathbb{R} \subseteq A_v$. Then v belongs to one of the two following types of real places:

- 1. For every $t \in \mathbb{R}$ there is a real place $v_t \colon \mathbb{R}(X) \to \mathbb{Z} \cup \{\infty\}$, defined by $v_t(y) = a$, where $a \in \mathbb{Z}$ is such that y can be written as $y = (X t)^a \cdot \frac{f}{g}$, with f, g being polynomials in $\mathbb{R}[X]$ not divisible by X t. The residue field $\mathbb{R}(X)_{v_t}$ is isomorphic to \mathbb{R} by the natural homomorphism $y + M_{v_t} \mapsto y(t)$.
- 2. There is a real place $v_{\infty} \colon \mathbb{R}(X) \to \mathbb{Z} \cup \{\infty\}$, defined by $v_t(\frac{f}{g}) = \deg g \deg f$, where f, g are polynomials in $\mathbb{R}[X]$. The residue field $\mathbb{R}(X)_{v_{\infty}}$ is isomorphic to \mathbb{R} by the natural homomorphism $y + M_{v_{\infty}} \mapsto \lim_{t \to \infty} y(t)$.

Let P be a semicone of a field F, and F_0 be a subfield of F. Denote by \leq the corresponding semiordering of P; define the set

$$A_{F_0}^P \coloneqq \{a \in F \mid a \le b \text{ and } -a \le b \text{ for some } b \in F_0\}.$$
(4)

The following lemmas show that $A_{F_0}^P$ is a valuation ring, and that its corresponding residue field admits a semiordering induced by P under additional conditions.

▶ Lemma 3.8 ([24, Lemma 7.13]). Let P be a semiordering of a field F and F_0 a subfield of F. Then $A_{F_0}^P$ is a valuation ring of some valuation of F.

▶ Lemma 3.9 ([24, Lemma 7.15]). Let P be a semiordering of a field F and F_0 a subfield of F, such that there exists $b \in F$ with $a \leq b$ for all $a \in F_0$. Let the valuation v of F correspond to $A_{F_0}^P$. Then $(A_v \cap P)/M_v$ is a semiordering of F_v .

4 Local-global principle over strictly positive polynomials

For a subset B of \mathbb{R} , define the set W(B) of polynomials that are non-negative on B:

 $W(B) \coloneqq \{ f \in \mathbb{R}[X] \mid f(x) \ge 0 \text{ for all } x \in B \}.$

Obviously $U(B) \subseteq W(B)$. For $f, g \in W(\mathbb{R}) \setminus \{0\}$, by the fundamental theorem of algebra, one can write (uniquely)

$$f = c \prod_{j \in J} (x - r_j)^{d_j} \prod_{k \in K} (x^2 + a_k x + b_k)^{e_k}, \quad g = c' \prod_{j \in J} (x - r_j)^{d'_j} \prod_{k \in K} (x^2 + a_k x + b_k)^{e'_k}$$

where $c, c', r_j, a_k, b_k \in \mathbb{R}$ and d_j, d'_j, e_k, e'_k are non-negative integers, and the polynomials $x^2 + a_k x + b_k$ have no real root. Here, J indexes all real roots of f and g, and K indexes all conjugate pairs of imaginary roots of f and g. Since f and g are non-negative on \mathbb{R} , all d_j and d'_j are even, and c, c' are positive. Therefore, the greatest common divisor of f and g, defined by

$$\gcd(f,g) \coloneqq \prod_{j \in J} (x - r_j)^{\min\{d_j, d'_j\}} \prod_{k \in K} (x^2 + a_k x + b_k)^{\min\{e_k, e'_k\}}$$

is also non-negative on \mathbb{R} . It follows that the polynomials gcd(f,g), f/gcd(f,g) and g/gcd(f,g) are all in $W(\mathbb{R})$.

We now give a proof of our first main result, which can be considered as a "strictly positive" version of the Bröcker-Prestel local-global principle. A comparison of our proof with the proof of the original theorem is given in Appendix C.

▶ **Theorem 2.1.** Let B be a subset of \mathbb{R} . Assume that $h_1, \ldots, h_n \in \mathbb{R}[X]$ are polynomials that satisfy gcd $(h_1, \ldots, h_n) = 1$. If the equation $f_1h_1 + \cdots + f_nh_n = 0$ has no solution (f_1, \ldots, f_n) over U(B), then there exists a real number $t \in \overline{B}$, such that the values $h_i(t), i = 1, \ldots, n$, are either all non-negative or all non-positive.

Proof. The theorem is trivially true if B is empty, hence we suppose $B \neq \emptyset$. Suppose $f_1h_1 + \cdots + f_nh_n = 0$ has no solution (f_1, \ldots, f_n) over U(B). Consider the following subset of the field $\mathbb{R}(X)$:

$$P_0 \coloneqq \left\{ \frac{g}{G} \cdot \sum_{i=1}^n f_i h_i, \text{ where all } f_i \in \mathcal{U}(B) \text{ and } g, G \in \mathcal{W}(\mathbb{R}) \setminus \{0\} \right\}$$

Since $f_1h_1 + \cdots + f_nh_n = 0$ has no solution (f_1, \ldots, f_n) over U(B), we have $0 \notin P_0$. We claim that $P'_0 = P_0 \cup \{0\}$ is a pre-semicone of $\mathbb{R}(X)$. Indeed, we verify the three conditions given in Definition 3.4:

(i) $P'_0 + P'_0 \subseteq P'_0$. It suffices to show $P_0 + P_0 \subseteq P_0$. Let $c = \frac{g}{G} \cdot \sum_{i=1}^n f_i h_i$, $c' = \frac{g'}{G'} \cdot \sum_{i=1}^n f'_i h_i$ be elements of P_0 . Without loss of generality we can suppose gcd(g, G) = gcd(g', G') = 1. Write $d \coloneqq gcd(g, g')$, $D \coloneqq gcd(G, G')$, then the polynomials $d, \frac{g}{d}, \frac{g'}{d}, D, \frac{G}{D}, \frac{G'}{D}$ are all elements of $W(\mathbb{R}) \setminus \{0\}$, and $gcd(\frac{gG'}{dD}, \frac{g'G}{dD}) = 1$. Hence,

$$c + c' = \sum_{i=1}^{n} \left(f_i \frac{g}{G} + f'_i \frac{g'}{G'} \right) h_i = \frac{dD}{GG'} \sum_{i=1}^{n} \left(f_i \frac{gG'}{dD} + f'_i \frac{g'G}{dD} \right) h_i$$
(5)

For any $x \in B$, we have $\frac{gG'}{dD}(x) \ge 0$ and $\frac{g'G}{dD}(x) \ge 0$. Since $\gcd(\frac{gG'}{dD}, \frac{g'G}{dD}) = 1$, the two polynomials $\frac{gG'}{dD}, \frac{g'G}{dD}$ cannot both vanish at x. Therefore either $\frac{gG'}{dD}(x) > 0$ or $\frac{g'G}{dD}(x) > 0$. Because $f_i(x) > 0$ and $f'_i(x) > 0$, it follows that $\left(f_i \frac{gG'}{dD} + f'_i \frac{g'G}{dD}\right)(x) > 0$. So $f_i \frac{gG'}{dD} + f'_i \frac{g'G}{dD} \in U(B)$ and $c + c' \in P_0$.

- So $f_i \frac{gG'}{dD} + f'_i \frac{g'G}{dD} \in U(B)$, and $c + c' \in P_0$. (ii) $\mathbb{R}(X)^2 \cdot P'_0 \subseteq P'_0$. This is obvious since $\mathbb{R}[X]^2 \cdot W(\mathbb{R}) \subseteq W(\mathbb{R})$.
- (iii) $P'_0 \cap -P'_0 = \{0\}$. It suffices to show $P_0 \cap -P_0 = \emptyset$. On the contrary suppose $c \in P_0 \cap -P_0$, then $0 = c + (-c) \in P_0 + P_0 \subseteq P_0$, a contradiction.

By Lemma 3.5, P'_0 can be extended to some P such that either P or -P is a semicone of the field $\mathbb{R}(X)$. Without loss of generality suppose $P \supseteq P'_0$ is a semicone, otherwise we can replace all h_i by $-h_i$. Since the field $\mathbb{R}(X)$ has no archimedean ordering [25, Example 1.1.4(2)], the *semiordering* corresponding to P must be non-archimedean (otherwise by Lemma 3.6 it must be an archimedean ordering). Consider the subfield \mathbb{R} of $\mathbb{R}(X)$, by Lemma 3.8 the valuation ring $A^P_{\mathbb{R}}$ (as defined in (4)) corresponds to some valuation v of $\mathbb{R}(X)$. Since P is non-archimedean, there exists some $a \in \mathbb{R}(X)$ such that $a - r \in P$ for all $r \in \mathbb{R}$, hence $A^P_{\mathbb{R}} \neq \mathbb{R}(X)$. Also, Lemma 3.9 shows that the residue field F_v admits a semiordering $(P \cap A_v)/M_v$. By Lemma 3.3, F_v is formally real. Therefore, v is a non-trivial real place of $\mathbb{R}(X)$, and from the definition of $A^P_{\mathbb{R}}$ we have $\mathbb{R} \subseteq A^P_{\mathbb{R}} = A_v$.

26:8 Solving Homogeneous Linear Equations over Polynomial Semirings

Using the classification of real places of $\mathbb{R}(X)$ given in Proposition 3.7, consider the following three cases. Since F_v is isomorphic to \mathbb{R} , the semiordering $(P \cap A_v)/M_v$ corresponds to the only ordering on \mathbb{R} .

1. The real place v is equivalent to a place v_t for some $t \in \overline{B} \subseteq \mathbb{R}$. In this case $\mathbb{R}[X] \subseteq A_v$. We show that $h_i(t) \ge 0$ for all i. By symmetry it suffices to show $h_1(t) \ge 0$. For every $\varepsilon \in \mathbb{R}_{>0}$, we have $\varepsilon \in U(B)$, so $h_1 + \varepsilon (h_2 + \cdots + h_n) \in P_0 \subseteq P$. Since $h_1 + \varepsilon (h_2 + \cdots + h_n) \in \mathbb{R}[X] \subseteq A_v$, we have $h_1 + \varepsilon (h_2 + \cdots + h_n) \in P \cap A_v$, which gives

$$h_1 + \varepsilon (h_2 + \dots + h_n) + M_v \in (P \cap A_v) / M_v.$$
(6)

Since the residue field $\mathbb{R}(X)_v$ is isomorphic to \mathbb{R} by the natural homomorphism $y + M_v \mapsto y(t)$, Equation (6) yields

$$h_1(t) + \varepsilon (h_2(t) + \dots + h_n(t)) \ge 0.$$

Since this is true for all $\varepsilon > 0$, we conclude that $h_1(t) \ge 0$ and thus $h_i(t) \ge 0$ for all *i*.

- 2. The real place v is equivalent to a place v_t for some $t \in \mathbb{R} \setminus \overline{B}$. There exists a polynomial $H_B \in \mathbb{R}[X]$, such that $H_B(x) > 0$ for all $x \in B$ but $H_B(t) < 0$. Indeed, since $t \notin \overline{B}$, there exists an interval $(t \delta, t + \delta)$ disjoint from B; it then suffices to take $H_B := (X t)^2 \delta^2$. As in the previous case, we have $h_1(t) \ge 0$. Furthermore, since $H_B \in U(B)$ by its definition, we have $H_Bh_1 + \varepsilon(h_2 + \cdots + h_n) \in P_0 \subseteq P$ for all $\varepsilon \in \mathbb{R}_{>0}$. This yields $(H_Bh_1)(t) \ge 0$. However, we have $H_B(t) < 0$ by its definition. This together with $h_1(t) \ge 0$ yields $h_1(t) = 0$. By symmetry we can prove $h_i(t) = 0$ for all i, this contradicts the condition $\gcd(h_1, \ldots, h_n) = 1$.
- 3. The real place v is equivalent to the place v_{∞} . We divide $\{h_1, \ldots, h_n\}$ into two parts according to the parity of its degree. Without loss of generality, suppose h_1, \ldots, h_k have even degree, and h_{k+1}, \ldots, h_n have odd degree.

Define the *leading coefficient* of a polynomial as the coefficient of its highest degree monomial. First we claim that the leading coefficients of h_1, \ldots, h_k are all positive. By symmetry, we only prove positivity of the leading coefficients of h_1 .

Let $m = \max\{\deg h_1, \ldots, \deg h_n\} + 1$. Since $(X^2 + 1)^m \in U(B)$ and $X^{\deg h_1} \in W(\mathbb{R})$, we have

$$\frac{h_1}{X^{\deg h_1}} + \frac{(X^2 + 1)^m}{(X^2 + 1)^{2m}}(h_2 + \dots + h_n) + M_v \in (P \cap A_v)/M_v.$$
(7)

Since the residue field $\mathbb{R}(X)_v$ is isomorphic to \mathbb{R} by the natural homomorphism $y + M_{v_t} \mapsto \lim_{t\to\infty} y(t)$, Equation (7) shows that the leading coefficient of h_1 is positive. Therefore by symmetry, the leading coefficient of h_i is positive for all $1 \leq i \leq k$. We then separate four cases.

a. If B is bounded, that is, $B \subset (a, b)$ for some $a, b \in \mathbb{R}$. Let $s > \max\{|a|, |b|\}$, then $X + s \in U(B)$. Since deg h_n is odd, we have $X^{\deg h_n + 1} \in W(\mathbb{R})$. Therefore,

$$\frac{(X+s)h_n}{X^{\deg h_n+1}} + \frac{(X^2+1)^m}{(X^2+1)^{2m}}(h_1+\dots+h_{n-1}) + M_v \in (P \cap A_v)/M_v.$$
(8)

This shows that the leading coefficient of h_n is positive.

However, we also have $-X + s \in U(B)$, so we can replace (X + s) with (-X + s) in Equation (8). This shows that the leading coefficient of h_n is negative. Therefore h_n does not exist, so all h_1, \ldots, h_n must have even degree. But then $(X+1-a)(b+1-X) \in U(B)$, so

$$\frac{(X+1-a)(b+1-X)h_1}{X^{\deg h_1+2}} + \frac{(X^2+1)^m}{(X^2+1)^{2m}}(h_2+\dots+h_n) + M_v \in (P \cap A_v)/M_v.$$
(9)

This shows that the leading coefficient of h_1 is negative, a contradiction.

b. If $B \subset (a, +\infty)$ for some $a \in \mathbb{R}$, and B contains arbitrarily large positive reals, that is, $B \cap (b, +\infty) \neq \emptyset$ for all $b \in \mathbb{R}$. Then $X + 1 - a \in U(B)$, so

$$\frac{(X+1-a)h_n}{X^{\deg h_n+1}} + \frac{(X^2+1)^m}{(X^2+1)^{2m}}(h_1+\dots+h_{n-1}) + M_v \in (P \cap A_v)/M_v.$$
(10)

This shows that the leading coefficient of h_n is positive. By symmetry, the leading coefficients of h_{k+1}, \ldots, h_n are all positive. Therefore, for large enough $t \in B$, $h_1(t), \ldots, h_n(t)$ are all positive.

c. If $B \subset (-\infty, a)$ for some $a \in \mathbb{R}$, and B contains arbitrarily small reals, that is, $B \cap (-\infty, b) \neq \emptyset$ for all $b \in \mathbb{R}$. Then $a + 1 - X \in U(B)$, so

$$\frac{(a+1-X)h_n}{X^{\deg h_n+1}} + \frac{(X^2+1)^m}{(X^2+1)^{2m}}(h_1+\dots+h_{n-1}) + M_v \in (P \cap A_v)/M_v.$$
(11)

This shows that the leading coefficient of h_n is negative. By symmetry, the leading coefficients of h_{k+1}, \ldots, h_n are all negative. Therefore, for small enough $0 > t \in B$, $h_1(t), \ldots, h_n(t)$ are all positive.

d. If *B* contains arbitrarily large and arbitrarily small reals. We claim that the leading coefficients of h_{k+1}, \ldots, h_n all have the same sign. Suppose on the contrary that they have different signs, denote by a_i the leading coefficient of h_i , so $h_i = a_i X^{\deg h_i} + H_i$ for some polynomial H_i of degree at most deg $h_i - 1$. Then there exist strictly positive reals r_{k+1}, \ldots, r_n such that $r_{k+1}a_{k+1} + \cdots + r_na_n = 0$. Then, for any $s \in \mathbb{R}$, we have $X^2 - 2sX + s^2 + 1 \in U(B)$, so

$$\frac{(X^2+1)^m}{(X^2+1)^{2m}}(h_1+\dots+h_k) + \frac{r_{k+1}(X^2-2sX+s^2+1)}{X^{\deg h_{k+1}+1}}h_{k+1} + \frac{r_{k+2}}{X^{\deg h_{k+2}-1}}h_{k+2} + \dots + \frac{r_n}{X^{\deg h_n-1}}h_n + M_v \in (P \cap A_v)/M_v.$$
(12)

The limit of the left hand side when X tends to infinity is equal to

$$g(s) \coloneqq \lim_{X \to \infty} \left(\frac{r_{k+1}(X^2 - 2sX)}{X^{\deg h_{k+1} + 1}} h_{k+1}(X) + \sum_{j=k+2}^n \frac{r_j}{X^{\deg h_j - 1}} h_j(X) \right)$$
$$= -2sa_{k+1}r_{k+1} + \lim_{X \to \infty} \left(\sum_{j=k+1}^n \frac{r_j}{X^{\deg h_j - 1}} H_j(X) \right)$$

because $r_{k+1}a_{k+1} + \cdots + r_na_n = 0$. According to whether $a_{k+1}r_{k+1}$ is positive or negative, we can take a positive or negative s with large enough absolute value, so that the value of g(s) is negative. This contradicts Equation (12), which shows that the limit of the left hand side when $X \to \infty$ is positive.

We therefore conclude that the leading coefficients of h_{k+1}, \ldots, h_n all have the same sign. If they are positive, then for large enough $t \in B$, $h_1(t), \ldots, h_n(t)$ are all positive. If they are negative, then for small enough $t \in B$, $h_1(t), \ldots, h_n(t)$ are all positive.

To sum up, in all possible cases, we have $t \in \overline{B}$ with $h_i(t) \ge 0$ for all *i*. If -P is a semicone instead of P, analogously we can find $t \in \overline{B}$ such that $h_i(t) \le 0$ for all *i*.

5 Local-global principle over $\mathbb{N}[X]$

In this section we prove Theorem 2.2. Omitted proofs are given in Appendix A. The key to bridging the difference between the semirings U(B) and $\mathbb{N}[X]$ is Pólya's Theorem:

▶ Lemma 5.1 (Pólya's Theorem [14, Theorem 56]). If a homogeneous polynomial $f \in \mathbb{R}[X_1, \ldots, X_n]$ is strictly positive for all (X_1, \ldots, X_n) on $(\mathbb{R}_{\geq 0})^n \setminus \{0\}$, then there exists $p \in \mathbb{N}$ such that $(X_1 + \cdots + X_n)^p \cdot f \in \mathbb{R}_{\geq 0}[X_1, \ldots, X_n]$.

The following proposition reduces Theorem 2.2 to real polynomials.

▶ **Proposition 5.2.** Given $h_1, \ldots, h_n \in \mathbb{Z}[X]$. The equation $f_1h_1 + \cdots + f_nh_n = 0$ has a solution (f_1, \ldots, f_n) over $\mathbb{N}[X] \setminus \{0\}$ if and only if it has a solution over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$.

The next proposition further reduces it to $U(\mathbb{R}_{>0})$. The key to its proof is Lemma 5.1.

▶ Proposition 5.3. Given $h_1, \ldots, h_n \in \mathbb{Z}[X]$. The equation $f_1h_1 + \cdots + f_nh_n = 0$ has a solution (f_1, \ldots, f_n) over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$ if and only if it has a solution over $U(\mathbb{R}_{>0})$.

This justifies the need for a "strictly positive" version of the Bröcker-Prestel principle, since Proposition 5.3 no longer holds if we replace $U(\mathbb{R}_{>0})$ with $W(\mathbb{R}_{>0}) \setminus \{0\}$ (see Remark A.1). We now prove the local-global principle for homogeneous linear equations over $\mathbb{N}[X]$.

▶ **Theorem 2.2.** Given polynomials $h_1, \ldots, h_n \in \mathbb{Z}[X]$ with $gcd(h_1, \ldots, h_n) = 1$. If the equation $f_1h_1 + \cdots + f_nh_n = 0$ has no solution (f_1, \ldots, f_n) over $\mathbb{N}[X] \setminus \{0\}$, then there exists $t \in \mathbb{R}_{>0}$, such that the values $h_i(t), i = 1, \ldots, n$ are either all non-negative or all non-positive.

Proof. Suppose the equation $f_1h_1 + \cdots + f_nh_n = 0$ has no solution (f_1, \ldots, f_n) over $\mathbb{N}[X] \setminus \{0\}$. By Proposition 5.2 and 5.3, it has no solution over $U(\mathbb{R}_{>0})$. Hence, by Theorem 2.1, there exists a real number $t \in \overline{\mathbb{R}_{>0}} = \mathbb{R}_{>0}$ such that $h_i(t)$ are all non-negative or all non-positive.

6 Decidability

In this section we show our main decidability result.

▶ **Theorem 2.3.** Given as input $h_1, \ldots, h_n \in \mathbb{Z}[X]$. It is decidable in polynomial time whether the equation $f_1h_1 + \cdots + f_nh_n = 0$ has solutions f_1, \ldots, f_n over $\mathbb{N}[X] \setminus \{0\}$.

Proof. (A summary of the algorithm constructed in this proof is given in Appendix B.)

By the homogeneity of the linear equation, we can divide h_1, \ldots, h_n by their greatest common divisor and suppose $gcd(h_1, \ldots, h_n) = 1$. Computing the greatest common divisor can be done in polynomial time using the Euclidean algorithm.

We then show that we can simplify the equation so that h_1, \ldots, h_n satisfy

$$h_i(0) > 0, h_j(0) < 0, \quad \text{for some } i, j.$$
 (13)

Suppose this is not already the case, that $h_i(0) \ge 0$ for all i or $h_i(0) \le 0$ for all i. Without loss of generality suppose $h_i(0) \ge 0$ for all i. We write $h_1(0) = 0, \ldots, h_k(0) = 0, h_{k+1}(0) > 0, \ldots, h_n(0) > 0$. Then $X \mid h_i$ for $i = 1, \ldots, k$.

If k = 0, that is $h_i(0) > 0$ for all i, then $f_1h_1 + \cdots + f_nh_n = 0$ has no solution over $\mathbb{N}[X] \setminus \{0\}$. Indeed, suppose on the contrary that (f_1, \ldots, f_n) is such a solution. Dividing all f_i by a suitable power of X we can suppose $f_s(0) \neq 0$ for some s. Then $f_i(0) \geq 0$ for all i while $f_s(0) > 0$, which yields $f_1(0)h_1(0) + \cdots + f_n(0)h_n(0) > 0$, a contradiction.

If $k \geq 1$, we show that the equation

$$f_1h_1 + \dots + f_nh_n = 0 \tag{14}$$

has a solution over $\mathbb{N}[X] \setminus \{0\}$ if and only if the equation

$$f_1 \cdot \frac{h_1}{X} + \dots + f_k \cdot \frac{h_k}{X} + f_{k+1}h_{k+1} + \dots + f_nh_n = 0$$
(15)

has a solution over $\mathbb{N}[X] \setminus \{0\}$. Let (f_1, \ldots, f_n) be a solution over $\mathbb{N}[X] \setminus \{0\}$ of Equation (14), then $f_1(0)h_1(0) + \cdots + f_n(0)h_n(0) = 0$. Since $h_i(0) = 0$ for all $i = 1, \ldots, k$, $h_i(0) > 0$ for $i = k + 1, \ldots, n$ and $f_i(0) \ge 0$ for $i = 1, \ldots, n$, we must have $f_{k+1}(0) = 0, \ldots, f_n(0) = 0$. That is, $X \mid f_{k+1}, \ldots, X \mid f_n$. Therefore $(f_1, \ldots, f_k, f_{k+1}/X, \ldots, f_n/X)$ is a solution over $\mathbb{N}[X] \setminus \{0\}$ of Equation (15). This shows that we can divide h_1, \ldots, h_k by X without changing the existence of solutions of Equation (14). Repeating this division process, one eventually terminates by obtaining h_i such that either: $h_i(0)$ are all strictly positive or all strictly negative, in which case Equation (14) has no solution over $\mathbb{N}[X] \setminus \{0\}$; or $h_i(0) > 0$ and $h_j(0) < 0$ for some i, j, in which case we have achieved the desired simplification to Condition (13). This procedure is repeated at most deg $h_1 + \cdots + \deg h_n$ times, and therefore terminates in polynomial time.

Supposing Condition (13), we claim that $f_1h_1 + \cdots + f_nh_n = 0$ has no solution over $\mathbb{N}[X] \setminus \{0\}$ if and only if there exists $t \in \mathbb{R}_{\geq 0}$ such that $h_i(t)$ are all non-positive or all non-negative. The first implication is given by Theorem 2.2. Conversely, suppose $h_i(t)$ are all non-positive or all non-negative. Without loss of generality suppose $h_i(t) \geq 0$ for all *i*. By Condition (13), we have $t \neq 0$. Suppose on the contrary that (f_1, \ldots, f_n) is a solution over $\mathbb{N}[X] \setminus \{0\}$, then $f_i(t) > 0$ for all *i* since t > 0. Since $gcd(h_1, \ldots, h_n) = 1$, at least one of $h_i(t)$ must be non-zero. Since $h_i(t) \geq 0$ for all *i*, we have $f_1(t)h_1(t) + \cdots + f_n(t)h_n(t) > 0$, a contradiction.

Thus, it suffices to decide whether there exists $t \ge 0$ such that $h_i(t)$ are all non-positive or all non-negative. This can be expressed in the existential theory of the reals:

$$\exists X \ (X \ge 0 \land h_1(X) \ge 0 \land \dots \land h_n(X) \ge 0) \lor (X \ge 0 \land h_1(X) \le 0 \land \dots \land h_n(X) \le 0).$$
(16)

Deciding the existential theory of the reals *in one variable* can be done in polynomial time with respect to the total bit length used to encode the sentence, due to a classic result by Collins^2 [8]. Therefore, one can decide the correctness of the sentence (16) in polynomial time. Combining all the steps, we conclude that the total complexity is in PTIME.

7 Application to wreath product

In this section we show the following result on wreath products.

▶ **Theorem 2.4.** Given a finite set of elements $\mathcal{G} = \{(y_1, b_1), \dots, (y_n, b_n)\}$ in $\mathbb{Z} \wr \mathbb{Z}$, where $b_i = \pm 1$ for all *i*. The following are decidable:

- **1.** (Group Problem) whether the semigroup $\langle \mathcal{G} \rangle$ generated by \mathcal{G} is a group.
- **2.** (Identity Problem) whether the neutral element I is in the semigroup $\langle \mathcal{G} \rangle$.

Let φ be the isomorphism defined in (3). Fix a finite set of elements \mathcal{G} as in Theorem 2.4. For $i = 1, \ldots, n$, denote by $H_i \in \mathbb{Z}[X, X^{-1}]$ the Laurent polynomial in the upper-right entry of the image of $\varphi((y_i, b_i))$. Write $\mathcal{G} = \mathcal{G}_+ \cup \mathcal{G}_-$ where $\mathcal{G}_+ := \{(y_i, b_i) \in \mathcal{G} \mid b_i = 1\}$ and $\mathcal{G}_- := \{(y_j, b_j) \in \mathcal{G} \mid b_j = -1\}$. Let $\varphi(\mathcal{G}), \varphi(\mathcal{G}_+), \varphi(\mathcal{G}_-)$ be the set of matrices that are images under φ of elements in $\mathcal{G}, \mathcal{G}_+, \mathcal{G}_-$. Define the sets of indices

 $I \coloneqq \{i \mid b_i = 1\}, \quad J \coloneqq \{i \mid b_i = -1\}.$

² The algorithm by Collins [8] has complexity $L^3(nd)^{2^{O(K)}}$, where L is the total coefficient bit length, n the number of polynomials, d the total degree of the polynomials, and K the number of variables. In the one variable case, K = 1, the algorithm takes polynomial time with respect to the total bit length.

For simplicity, we write $A_i, i \in I$ for the matrices in $\varphi(\mathcal{G}_+)$, and $B_j, j \in J$ the matrices in $\varphi(\mathcal{G}_-)$. For every tuple $(i, j) \in I \times J$, define the Laurent polynomial

$$h_{ij} \coloneqq X^{-1}H_i + H_j \in \mathbb{Z}[X, X^{-1}].$$

$$\tag{17}$$

This is the upper-right entry of the matrix $A_i B_j$.

For a subset $S \subseteq I \times J$, denote by $\pi_I(S)$ its projection onto the I coordinates, that is, $\pi_I(S) := \{i \in I \mid \exists j \in J, (i, j) \in S\}$. Define $\pi_J(S)$ likewise. The key to proving the partial decidability of the Group Problem in $\mathbb{Z} \wr \mathbb{Z}$ is the following proposition that relates sub-semigroups of $\mathbb{Z} \wr \mathbb{Z}$ to equations over $\mathbb{N}[X] \setminus \{0\}$.

▶ **Proposition 7.1.** Given a set $\mathcal{G} = \mathcal{G}_+ \cup \mathcal{G}_-$ of generators defined as above. Let $h_{ij} \in \mathbb{Z}[X, X^{-1}]$ be the polynomials defined in (17). The semigroup $\langle \mathcal{G} \rangle$ is a group if and only if there exists a set $S \subseteq I \times J$ satisfying $\pi_I(S) = I, \pi_J(S) = J$, such that the equation $\sum_{(i,j)\in S} f_{ij}h_{ij} = 0$ has a solution $(f_{ij})_{(i,j)\in S}$ over $\mathbb{N}[X] \setminus \{0\}$.

Proof. For a word w in the alphabet $\varphi(\mathcal{G})$, define its product $\pi(w)$ to be the matrix obtained by multiplying all the matrices in w consecutively. Denote by $|w|_+$ (respectively, $|w|_-$) the number of letters in w belonging in $\varphi(\mathcal{G}_+)$ (respectively, $\varphi(\mathcal{G}_-)$). Define the *height* of the word w to be $h(w) := |w|_+ - |w|_-$, then we have $\pi(w) = \begin{pmatrix} 1 & * \\ 0 & X^{h(w)} \end{pmatrix}$, where * is some element in $\mathbb{Z}[X, X^{-1}]$.

For a finite alphabet \mathcal{A} , denote by \mathcal{A}^+ the set of non-empty words over \mathcal{A} . We claim that for any non-empty word $w \in \varphi(\mathcal{G})^+$ such that h(w) = 0, the upper right entry of $\pi(w)$ can be written as a sum $\sum_{(i,j)\in I\times J} f_{ij}h_{ij}$, where f_{ij} are elements in $\mathbb{N}[X, X^{-1}]$. We prove this by induction on the length of the word w. For the sake of simplicity, denote $U(\pi(w))$ the upper right entry of $\pi(w)$.

If w has length at most two, then it must be of the form A_iB_j or B_jA_i , and the claim is easy to verify. Suppose the claim is true for all words w of length less then $\ell > 2$. We prove the claim for words w of length ℓ . Distinguish the following two cases.

1. The word w is of the form $A_i w' B_j$ or $B_j w' A_i$ for some $i \in I, j \in J, w' \in \varphi(\mathcal{G})^+$. Since w' has length at most $\ell - 2$ and is of height 0, by induction hypothesis, $\pi(w') = \begin{pmatrix} 1 & r \\ 0 & 1 \end{pmatrix}$, with r a linear combination of h_{ij} with coefficients in $\mathbb{N}[X, X^{-1}]$. If $w = A_i w' B_j$, then

$$\pi(w) = \begin{pmatrix} 1 & H_i \\ 0 & X \end{pmatrix} \begin{pmatrix} 1 & r \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & H_j \\ 0 & X^{-1} \end{pmatrix} = \begin{pmatrix} 1 & X^{-1}r + (X^{-1}H_i + H_j) \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & X^{-1}r + h_{ij} \\ 0 & 1 \end{pmatrix}.$$

So $U(\pi(w)) = X^{-1}r + h_{ij}$ can also be written as a linear combination of $h_{ij}, i \in I, j \in J$ with coefficients in $\mathbb{N}[X, X^{-1}]$. If $w = B_j w' A_i$, then

$$\pi(w) = \begin{pmatrix} 1 & H_j \\ 0 & X^{-1} \end{pmatrix} \begin{pmatrix} 1 & r \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & H_i \\ 0 & X \end{pmatrix} = \begin{pmatrix} 1 & Xr + H_i + XH_j \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & X(r+h_{ij}) \\ 0 & 1 \end{pmatrix}$$

So $U(\pi(w)) = X(r + h_{ij})$ can also be written as a linear combination of $h_{ij}, i \in I, j \in J$ with coefficients in $\mathbb{N}[X, X^{-1}]$.

2. The word w is of the form $A_iw'A_{i'}$ or $B_jw'B_{j'}$ for some $i, i' \in I$ or $j, j' \in J$. First suppose $w = A_iw'A_{i'}$. Since $h(A_i) = 1 > 0$ and $h(A_iw') = -1 < 0$, there must exist a strict prefix v of w with height zero. This is because by reading the word wletter by letter, this height of consecutive prefixes differs by at most one. We have w = vv' with h(v) = h(v') = 0 where v, v' are non-empty words. By induction hypothesis,

 $U(\pi(v)), U(\pi(v'))$ can be written as a linear combination of h_{ii} with coefficients in $\mathbb{N}[X, X^{-1}]$. Therefore $U(\pi(w)) = U(\pi(v)) + U(\pi(v'))$ also satisfies this claim. The case where $w = B_i w' B_{i'}$ is completely analogous.

Combining the two cases concludes the induction. It is easy to see from the induction process that if the letter A_i appears in w, then the coefficient of the term h_{ij} in the linear combination is not zero for some $j \in J$. This is because at some point we have replaced r with either $X^{-1}r + h_{ij}$ or $X(r + h_{ij})$. Similarly, if the letter B_j appears in w, then the coefficient of the term h_{ij} in the linear combination is non-zero for some $i \in I$.

If the semigroup $\langle \mathcal{G} \rangle$ is a group, then there exists a word v in the alphabet \mathcal{G} using all letters in \mathcal{G} , whose corresponding product is the neutral element. Taking the image under φ yields a word $w = \varphi(v)$ in the alphabet $\varphi(\mathcal{G})$ such that h(w) = 0 and $U(\pi(w)) = 0$. The claim above and the discussion following it show that there exist Laurent polynomials $f_{ij} \in \mathbb{N}[X, X^{-1}]$ such that $\sum_{(i,j)\in I\times J} f_{ij}h_{ij} = 0$. Furthermore, all letters $A_i, i \in I$ and $B_j, j \in J$ appear in w, so for every i, the coefficient f_{ij} in the linear combination is not zero for some $j \in J$; and for every j, the coefficient f_{ij} is not zero for some $i \in I$. Let $S := \{(i,j) \in I \times J \mid f_{ij} \neq 0\}$, then $\sum_{(i,j) \in S} f_{ij}h_{ij} = 0$, and $\pi_I(S) = I, \pi_J(S) = J$. By the homogeneity of the equation $\sum_{(i,j)\in S} f_{ij}h_{ij} = 0$, one can multiply all f_{ij} by the monomial X^n for a sufficiently large n, and suppose $f_{ij} \in \mathbb{N}[X] \setminus \{0\}$ instead of $\mathbb{N}[X, X^{-1}] \setminus \{0\}$. This completes the proof of the first direction of implication in Proposition 7.1.

For the other direction of implication, suppose there exists a set $S \subseteq I \times J$ satisfying $\pi_I(S) = I, \pi_J(S) = J$, such that the equation $\sum_{(i,j)\in S} f_{ij}h_{ij} = 0$ has a solution $(f_{ij})_{(i,j)\in S}$ over $\mathbb{N}[X] \setminus \{0\}$. By the homogeneity of the equation, suppose that there is a tuple $(u, v) \in S$ such that $X \nmid f_{uv}$. Let $(y, z) \in S$ be a tuple such that $\deg f_{yz} \ge \deg f_{ij}$ for all $(i, j) \in S$. Denote by $\mathbb{N}_{>0}[X]$ the set of polynomials of the form $\sum_{i=0}^{d} a_i X^i$, where $d \ge 0$ and $a_i > 0$.

for all i. By multiplying all f_{ij} by the polynomial $(1 + X)^m$ for a sufficiently large m, we can suppose that $f_{uv} \in \mathbb{N}_{>0}[X], X^{-v_0(f_{yz})}f_{yz} \in \mathbb{N}_{>0}[X]$, and deg $f_{uv} \ge v_0(f_{yz})$. Indeed, we can take any $m \ge \max\{\deg f_{uv}, \deg X^{-v_0(f_{yz})}f_{yz}, v_0(f_{yz})\}$. Additionally, the condition that $\deg f_{yz} \geq \deg f_{ij}$ for all $(i,j) \in S$ is still satisfied after this multiplication.

We now construct a word $w \in \varphi(\mathcal{G})^+$ that uses every letter in $\varphi(\mathcal{G})$, such that $h(\pi(w)) = 0$, $U(\pi(w)) = \sum_{(i,j) \in S} f_{ij}h_{ij} = 0$. We start with the word

$$w_0 \coloneqq A_u^{\deg f_{uv}} A_y^{\deg f_{yz} - \deg f_{uv}} B_z^{\deg f_{yz} - \deg f_{uv}} B_v^{\deg f_{uv}},$$

which has height 0, and whose product has upper-right entry

$$U(\pi(w_0)) = h_{uv} \cdot \sum_{i=0}^{\deg f_{uv} - 1} X^i + h_{yz} \cdot \sum_{i=\deg f_{uv}}^{\deg f_{yz} - 1} X^i.$$

Since $f_{uv} \in \mathbb{N}_{>0}[X], X^{-v_0(f_{yz})}f_{yz} \in \mathbb{N}_{>0}[X]$, and deg $f_{uv} \ge v_0(f_{yz})$, the polynomials $\hat{f}_{uv} \coloneqq f_{uv} - \sum_{i=0}^{\deg f_{uv}-1} X^i, \ \hat{f}_{yz} \coloneqq f_{yz} - \sum_{i=\deg f_{uv}}^{\deg f_{yz}-1} X^i$ are still polynomials in $\mathbb{N}[X] \setminus \{0\}$. For $(i, j) \in S$, define

$$\hat{f}_{ij} \coloneqq \begin{cases} \hat{f}_{uv} & (i,j) = (u,v) \\ \hat{f}_{yz} & (i,j) = (y,z) \\ f_{ij} & \text{otherwise.} \end{cases}$$

These are elements in $\mathbb{N}[X] \setminus \{0\}$ and satisfy $U(\pi(w_0)) + \sum_{(i,j) \in S} \hat{f}_{ij} h_{ij} = \sum_{(i,j) \in S} f_{ij} h_{ij} = 0$. We then gradually insert "loops" of the form $A_i B_j$ into the word w_0 . This insertion does not change the height of the word, but it adds a multiple of h_{ij} to the upper-right entry of the product. Indeed, if h(vv') = 0, then we have $h(vA_iB_iv') = 0$ and $U(\pi(vA_iB_iv')) = 0$

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 $U(\pi(vv')) + X^{h(v)}h_{ij}$. Note that the initial word w_0 has suffixes of all heights from 0 to deg f_{yz} . For each $k = 0, \ldots$, deg f_{yz} and each $(i, j) \in S$, after a suffix of height k, we insert $\operatorname{Coef}_{X^k}(\hat{f}_{ij})$ times the "loop" A_iB_j , where $\operatorname{Coef}_{X^k}(\hat{f}_{ij})$ is the coefficient of the monomial X^k in the polynomial \hat{f}_{ij} . The upper-right entry of the product after all these insertions will be

$$U(\pi(w_0)) + \sum_{k=0}^{\deg f_{yz}} \sum_{(i,j)\in S} \operatorname{Coef}_{X^k}(\hat{f}_{ij}) X^k \cdot h_{ij} = U(\pi(w_0)) + \sum_{(i,j)\in S} \hat{f}_{ij} h_{ij} = 0,$$

because deg $f_{yz} \ge \text{deg } f_{ij}$ for all $(i, j) \in I \times J$. See Figure 1 for an example of this construction.

We have thus constructed a word $w \in \varphi(\mathcal{G})^+$ such that $h(\pi(w)) = 0$, $U(\pi(w)) = 0$. Note that we have inserted at least one loop A_iB_j for each $(i,j) \in S$. Since $\pi_I(S) = I$, $\pi_J(S) = J$, the word w contains every letter $A_i, i \in I$ and $B_j, j \in J$. Because $\pi(w)$ is the neutral element, the inverse of every letter in w can be written as a product of matrices in $\varphi(\mathcal{G})$. Indeed, if w = vXv' then $X^{-1} = \pi(v'v)$. Thus the inverse of every element of $\varphi(\mathcal{G})$ is in $\langle \varphi(\mathcal{G}) \rangle$. We conclude that $\langle \varphi(\mathcal{G}) \rangle$, and thus $\langle \mathcal{G} \rangle$, is a group.



Figure 1 Example of a word constructed in the proof of Proposition 7.1. Here, $S = \{(u,v), (y,z), (1,2), (3,1)\}$, and $f_{uv} = 1 + X + X^2 + X^3$, $f_{yz} = X^3 + X^4 + X^5 + X^6$, $f_{12} = X + 2X^5$, $f_{31} = 3 + X^2$. The constructed word is $A_u(A_1B_2)A_uA_u(A_uB_v)A_yA_y(A_1B_2)(A_1B_2)A_y(A_yB_z)B_zB_zB_zB_v(A_3B_1)B_vB_v(A_3B_1)(A_3B_1)(A_3B_1)$.

We have thus established the link between the Group Problem in $\mathbb{Z} \wr \mathbb{Z}$ and homogeneous linear equations over $\mathbb{N}[X]$. Theorem 2.4 follows from Proposition 7.1 and the decidability result of Theorem 2.3. Its proof is given in Appendix A.

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A Omitted proofs

▶ **Proposition 3.7.** Let v be a non-trivial real place of $\mathbb{R}(X)$ such that $\mathbb{R} \subseteq A_v$. Then v belongs to one of the two following types of real places:

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- 1. For every $t \in \mathbb{R}$ there is a real place $v_t \colon \mathbb{R}(X) \to \mathbb{Z} \cup \{\infty\}$, defined by $v_t(y) = a$, where $a \in \mathbb{Z}$ is such that y can be written as $y = (X t)^a \cdot \frac{f}{g}$, with f, g being polynomials in $\mathbb{R}[X]$ not divisible by X t. The residue field $\mathbb{R}(X)_{v_t}$ is isomorphic to \mathbb{R} by the natural homomorphism $y + M_{v_t} \mapsto y(t)$.
- 2. There is a real place $v_{\infty} \colon \mathbb{R}(X) \to \mathbb{Z} \cup \{\infty\}$, defined by $v_t(\frac{f}{g}) = \deg g \deg f$, where f, g are polynomials in $\mathbb{R}[X]$. The residue field $\mathbb{R}(X)_{v_{\infty}}$ is isomorphic to \mathbb{R} by the natural homomorphism $y + M_{v_{\infty}} \mapsto \lim_{t \to \infty} y(t)$.

Proof. Since $\mathbb{R} \subseteq A_v$, every element $r \in \mathbb{R} \setminus \{0\}$ satisfies $v(r) \ge 0$ and $v(r^{-1}) \ge 0$. But $v(r) + v(r^{-1}) = v(1) = 0$, so v(r) = 0. Consider the value v(X), there are two possibilities:

- 1. If $v(X) \ge 0$. In this case, we have $\mathbb{R} \subseteq A_v$ and $X \in A_v$, therefore $\mathbb{R}[X] \subseteq A_v$. Since M_v is a maximal (hence prime) ideal of A_v , the ideal $\mathbb{R}[X] \cap M_v$ is a prime ideal of $\mathbb{R}[X]$. Furthermore, $\mathbb{R}[X] \cap M_v$ is not zero, otherwise every element of $\mathbb{R}[X] \setminus \{0\}$ would be invertible in A_v , so $\mathbb{R}(X) \subseteq A_v$, contradicting the non-triviality of v. Since $\mathbb{R}[X]$ is a principle ideal domain, the non-zero prime ideal $\mathbb{R}[X] \cap M_v$ is generated by a single irreducible polynomial in $\mathbb{R}[X]$. Consider the two cases:
 - a. The ideal $\mathbb{R}[X] \cap M_v$ is generated by a polynomial X t for some $t \in \mathbb{R}$. In this case we have v(X t) > 0. Every polynomial $f \in \mathbb{R}[X]$ not divisible by (X t) can be written as $f = (X t) \cdot F + r$ for some $F \in \mathbb{R}[X]$, $r \in \mathbb{R} \setminus \{0\}$. Since $v((X t) \cdot F) = v(X t) + v(F) > 0$ and v(r) = 0, we have v(f) = v(r) = 0.

Every element $y \in \mathbb{R}(X)$ can be written as $y = (X-t)^a \cdot \frac{f}{g}$, where f, g are polynomials in $\mathbb{R}[X]$ not divisible by (X-t). Then $v(y) = a \cdot v(X-t) + v(f) - v(g) = av(X-t)$. Under isomorphism of the value group Γ , we can without loss of generality suppose v(X-t) = 1, then we get the valuation v_t of type 1 described in the proposition. Since every element $y \in M_{v_t}$ satisfies y(t) = 0, we have that $y + M_{v_t} \mapsto y(t)$ is an isomorphism from the residue field to \mathbb{R} ; it is a formally real field.

- **b.** The ideal $\mathbb{R}[X] \cap M_v$ is generated by a polynomial $X^2 + cX + d$ without real roots. In this case, the residue field A_v/M_v is a quadratic extension of \mathbb{R} , and is hence isomorphic to the field \mathbb{C} . However \mathbb{C} is not formally real. Indeed, suppose on the contrary that \mathbb{C} admits some ordering \leq , then since $0 < i^2 = -1$ and $0 < 1^2 = 1$, we have 0 < (-1) + 1 = 0, a contradiction.
- 2. If v(X) < 0. In this case we have $\mathbb{R}[1/X] \subseteq A_v$ and $1/X \in M_v$. Since $\mathbb{R}[1/X] \cap M_v$ is a prime ideal of $\mathbb{R}[1/X]$ that contains 1/X, it is generated by 1/X. Then similar to the case 1.a., every element $y \in \mathbb{R}(X)$ can be written as $y = (1/X)^a \cdot \frac{F}{G}$, where F, G are polynomials in $\mathbb{R}[1/X]$ not divisible by 1/X. Without loss of generality suppose v(1/X) = 1, we have v(y) = a. Rewrite $y = \frac{f}{g}$, comparing degrees, we have $a = \deg g \deg f$. So v is the valuation v_{∞} of type 2 described in the proposition. Since every element $y \in M_{v_{\infty}}$ satisfies $\lim_{t \to \infty} y(t) = 0$, we have that $y + M_{v_{\infty}} \mapsto \lim_{t \to \infty} y(t)$ is an isomorphism from the residue field to \mathbb{R} ; it is a formally real field.

▶ Proposition 5.2. Given $h_1, \ldots, h_n \in \mathbb{Z}[X]$. The equation $f_1h_1 + \cdots + f_nh_n = 0$ has a solution (f_1, \ldots, f_n) over $\mathbb{N}[X] \setminus \{0\}$ if and only if it has a solution over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$.

Proof. A solution over $\mathbb{N}[X] \setminus \{0\}$ is obviously also a solution over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$. Conversely, let $f_i = \sum_{j=0}^{d_i} a_{ij}X^j$, $i = 1, \ldots, n$, be a solution of $f_1h_1 + \cdots + f_nh_n = 0$. Write $h_i = \sum_{j=0}^{e_i} b_{ij}X^j$, $i = 1, \ldots, n$, then the equation $f_1h_1 + \cdots + f_nh_n = 0$ is equivalent to the system of equations

$$\sum_{i=1}^{n} \sum_{j=0}^{d} a_{ij} b_{i,d-j} = 0, \quad d = 1, \dots, \max_{1 \le i \le n} (d_i + e_i).$$
(18)

All the coefficients b_{ij} are integers, and $b_{i,d-j} = 0$ whenever d - j < 0.

If $f_1h_1 + \cdots + f_nh_n = 0$ has a solution over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$, then System (18) has a solution $a_{ij}, i = 1, \ldots, n, j = 1, \ldots, d_i$ over \mathbb{R} , satisfying

$$a_{ij} \ge 0, \quad i = 1, \dots, n, \quad j = 1, \dots, d_i,$$
(19)

and

$$a_{i1} \neq 0 \text{ or } a_{i2} \neq 0 \text{ or } \dots \text{ or } a_{id_i} \neq 0, \quad i = 1, \dots, n.$$
 (20)

This condition is a boolean combination of homogeneous linear inequalities with integer coefficients. Since the linear Systems (18), (19) and (20) have only integer coefficients, they have a solution over \mathbb{R} if and only if they have a solution over \mathbb{Q} . Then, by their homogeneity, they have a solution over \mathbb{Q} if and only if they have a solution over \mathbb{Z} . Hence, the Systems (18), (19), (20) have a solution over \mathbb{Z} , meaning $f_1h_1 + \cdots + f_nh_n = 0$ has a solution $f_i = \sum_{j=0}^{d_i} a_{ij}X^j$, $i = 1, \ldots, n$, over $\mathbb{N}[X] \setminus \{0\}$.

▶ Proposition 5.3. Given $h_1, \ldots, h_n \in \mathbb{Z}[X]$. The equation $f_1h_1 + \cdots + f_nh_n = 0$ has a solution (f_1, \ldots, f_n) over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$ if and only if it has a solution over $U(\mathbb{R}_{>0})$.

Proof. Obviously a solution over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$ is a solution over $U(\mathbb{R}_{\geq 0})$.

For the other implication, we use Pólya's Theorem (Lemma 5.1). Suppose $f_1h_1 + \cdots + f_nh_n = 0$ has a solution (f_1, \ldots, f_n) over $U(\mathbb{R}_{>0})$. Write $f_i = X^{c_i} \cdot F_i$ where $c_i \ge 0$ and $F_i \in \mathbb{R}[X]$ is such that $X \nmid F_i$. Since $X \nmid F_i$ we have $F_i(0) \ne 0$, we claim that $F_i(0) > 0$. In fact, if $F_i(0) < 0$, then by the continuity of F_i , there exists $\varepsilon > 0$ such that $F_i(\varepsilon) < 0$, but then $f_i(\varepsilon) = \varepsilon^{c_i} F_i(\varepsilon) < 0$, contradicting the fact that $f_i \in U(\mathbb{R}_{>0})$. Furthermore, one easily sees that $F_i(x) = \frac{f_i(x)}{x^{c_i}} > 0$ for all x > 0. So we have shown $F_i(x) > 0$ for all $x \ge 0$.

We now show that for large enough $p \in \mathbb{N}$, the polynomials $\hat{f}_i \coloneqq (X+1)^p \cdot f_i$ are all in $\mathbb{R}_{\geq 0}[X]$. Let Y be a new variable, and for every *i*, let G_i be the homogenization of F_i using the variable Y. That is, $G_i = F_i(X/Y) \cdot Y^{\deg(F_i)}$. Since $F_i(x/y) > 0$ for all $x/y \ge 0$, we have $G_i(x, y) > 0$ for all $x \ge 0, y > 0$. Whereas for $x > 0, y = 0, G_i(x, y)/x^{\deg(F_i)}$ is the leading coefficient of F_i . This is non-zero and thus must be positive because $\lim_{x\to\infty} F_i(x) > 0$. Therefore $G_i(x, y) > 0$ for x > 0, y = 0.

We have thus shown $G_i(x, y) > 0$ for all $x \ge 0, y \ge 0, x + y > 0$. Applying Pólya's Theorem yields the existence of a $p_i \in \mathbb{N}$ such that $(X+Y)^{p_i} \cdot G_i \in \mathbb{R}_{\ge 0}[X, Y]$. Taking Y = 1 we dehomogenize G_i and obtain $(X+1)^{p_i} \cdot F_i \in \mathbb{R}_{\ge 0}[X]$. Let $p = \max\{p_1, \ldots, p_n\}$, then

$$\hat{f}_i = (X+1)^p \cdot f_i = X^{c_i} \cdot (X+1)^p \cdot F_i \in \mathbb{R}_{\geq 0}[X] \setminus \{0\}$$

for all *i*. We have thus found the solution $(\hat{f}_1, \ldots, \hat{f}_n)$ over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$ for the equation $f_1h_1 + \cdots + f_nh_n = 0$.

▶ Remark A.1. Proposition 5.3 no longer holds if we replace $U(\mathbb{R}_{>0})$ with $W(\mathbb{R}_{>0}) \setminus \{0\}$. For example, take $n = 2, h_1 = 1, h_2 = -(X-1)^2$. Then $f_1 = (X-1)^2, f_2 = 1$ is a solution over $W(\mathbb{R}_{>0}) \setminus \{0\}$ of the equation $f_1h_1 + f_2h_2 = 0$. However, $f_1h_1 + f_2h_2 = 0$ does not admit a solution over $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$. Indeed, any solution of $f_1 - f_2 \cdot (X-1)^2 = 0$ over $\mathbb{R}[X]$ must satisfy $f_1(1) = 0$, so f_1 cannot be in $\mathbb{R}_{\geq 0}[X] \setminus \{0\}$.

▶ **Theorem 2.4.** Given a finite set of elements $\mathcal{G} = \{(y_1, b_1), \ldots, (y_n, b_n)\}$ in $\mathbb{Z} \wr \mathbb{Z}$, where $b_i = \pm 1$ for all *i*. The following are decidable:

1. (Group Problem) whether the semigroup $\langle \mathcal{G} \rangle$ generated by \mathcal{G} is a group.

2. (Identity Problem) whether the neutral element I is in the semigroup $\langle \mathcal{G} \rangle$.

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Proof.

- 1. For the Group Problem, by Proposition 7.1 it suffices to decide whether there exists a set $S \subseteq I \times J$ satisfying $\pi_I(S) = I, \pi_J(S) = J$, such that the equation $\sum_{(i,j)\in S} f_{ij}h_{ij} = 0$ has a solution $(f_{ij})_{(i,j)\in S}$ over $\mathbb{N}[X] \setminus \{0\}$. By the homogeneity of the equation $\sum_{(i,j)\in S} f_{ij}h_{ij} = 0$, one can multiply all the Laurent polynomials h_{ij} by a power of X and suppose all $h_{ij} \in \mathbb{Z}[X]$. For every set $S \subseteq I \times J$ satisfying $\pi_I(S) = I, \pi_J(S) = J$, we can use Theorem 2.3 to decide whether $\sum_{(i,j)\in S} f_{ij}h_{ij} = 0$ has a solution over $\mathbb{N}[X] \setminus \{0\}$. This shows the decidability of the Group Problem.
- 2. The neutral element is in $\langle \mathcal{G} \rangle$ if and only if a non-empty subset of \mathcal{G} generates a group (as a semigroup). This is because, if the product of a word $w \in \mathcal{G}^+$ is the neutral element, then every element in the set C of letters used in w can be inverted in $\langle C \rangle$, so $\langle C \rangle$ is a group. Therefore, in order to decide whether the neutral element is in $\langle \mathcal{G} \rangle$, it suffices to check for all subsets of \mathcal{G} whether they generate a group. This is decidable by the above result on the Group Problem.

B Algorithm for Theorem 2.2

Algorithm 1 Deciding existence of solutions over $\mathbb{N}[X] \setminus \{0\}$ of the equation $f_1h_1 + \cdots + f_nh_n = 0$.

Input: Polynomials $h_1, \ldots, h_n \in \mathbb{Z}[X]$. **Output: True** or **False**.

- (1) Compute $d \coloneqq \operatorname{gcd}(h_1, \ldots, h_n)$ and divide all h_i by d.
- (2) Repeat the following:
 - **a.** If $h_i(0) > 0$ for all i, or $h_i(0) < 0$ for all i, return **False**.
 - **b.** Else if $h_i(0) \ge 0$ for all i, or $h_i(0) \le 0$ for all i, divide all the polynomials h_i that satisfy $h_i(0) = 0$ by X.
 - c. Else go to 3.
- (3) Decide the truth of the existential sentence (16) in the theory of reals. If (16) is true, return **False**, otherwise return **True**.

C Comparison with the Bröcker-Prestel local-global principle

The original Bröcker-Prestel local-global principle ([24, Theorem 8.13]) can be formulated as follows.

▶ **Theorem C.1** (Bröcker-Prestel local-global principle). Let *F* be a formally real field, and h_1, \ldots, h_n be non-zero elements of *F*. If the equation $f_1h_1 + \cdots + f_nh_n = 0$ has no non-trivial solution $(f_1, \ldots, f_n) \neq (0, \ldots, 0)$ over sums of squares of *F* (that is, over the set $S := \{\sum_{i=1}^k a_i^2 \mid a_i \in F\}$), then at least one of the following hold:

(i) h_1, \ldots, h_n are all of the same sign in some archimedean ordering of F.

(ii) $f_1h_1 + \cdots + f_nh_n = 0$ has no solution in the Henselization of some real place of F.

For a definition of Henselizations of a formally real field, see [24, Proposition 8.1].

When applied to the field $F = \mathbb{R}(X)$, the Bröcker-Prestel local-global principle characterizes the absence of non-trivial solutions over sums of squares by condition (ii), since the field $\mathbb{R}(X)$ has no archimedean orderings. Multiplying by the common denominator and using the fact that any element in W(\mathbb{R}) can be written as a sum of squares in $\mathbb{R}(X)$,

Theorem C.1 also characterizes the absence of non-trivial solutions over $W(\mathbb{R})$. However, when considering non-trivial solutions over $U(\mathbb{R})$ and U(B), the situation is quite different; and we now compare the proof of Theorem 2.1 to Theorem C.1.

The proof of Bröcker-Prestel's original theorem starts with the definition of the presemicone

$$P_1 := \left\{ \sum_{i=1}^n f_i h_i, \text{ where } f_i \text{ are sum of squares of elements in } \mathbb{R}(X) \right\}.$$

Since it considers solutions over sum of squares, this definition is straightforward. The definition of P_0 is our proof of Theorem 2.1 is different and less straightforward. In our theorem, we are considering *strictly* positive polynomials on $B \subseteq \mathbb{R}$, therefore we need to replace sum of squares with polynomials in U(B). However, such a naive replacement does not work due to the requirement of a pre-semicone to be closed under multiplication of squares (unlike $W(\mathbb{R})$, the set U(B) is not closed under multiplication by squares). This is why we need to add the rational function $\frac{q}{G}$ in the definition of P_0 and use the fundamental theorem of algebra to guarantee closure under addition.

Note that in order to guarantee the closure under addition of P_0 , it is essential that we work in the *univariate* polynomial ring $\mathbb{R}[X]$, so that two polynomials g, g' having a common root implies $gcd(g, g') \neq 1$. For example, this no longer holds in the bivariate polynomial ring $\mathbb{R}[X, Y]$. Therefore, even when supposing $gcd(\frac{gG'}{dD}, \frac{g'G}{dD}) = 1$, we no longer have $\left(f_i \frac{gG'}{dD} + f'_i \frac{g'G}{dD}\right)(x, y) > 0$ in Equation (5). Thus, for the field $\mathbb{R}(X, Y)$, the closure under addition of P_0 no longer holds, a contrast with the "non-strict" version P_1 .

The following step of extracting the valuation ring $A_{\mathbb{R}}^{P}$ from the semiordering P appeared as part of the proof of the original theorem. (The original theorem used the valuation ring $A_{\mathbb{Q}}^{P}$ instead, but they are in fact equivalent.) This is the main part where we drew inspiration from the original local-global principle.

After extracting the valuation ring $A_{\mathbb{R}}^P$, our proof again diverges from that of the original theorem. Our new definition of P_0 allows us to enforce strict positivity, however it also takes away some convenient properties of the pre-semicone P_1 in the original theorem. Notably, we have $h_1 \in P_1$, allowing for a quick conclusion on the positivity of h_1 in Henselizations. Whereas for P_0 , we do not have $h_1 \in P_0$ due to the strict positivity of the coefficients f_i . We compensate this by the analytic approach adopted in the second half of our proof, making use of the classification of real places of $\mathbb{R}(X)$ and the continuity of functions in $\mathbb{R}[X]$. This part is absent from the proof of the original theorem, which is purely algebraic and model theoretic.