# Reconfiguration and Locomotion with Joint Movements in the Amoebot Model 

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#### Abstract

We are considering the geometric amoebot model where a set of $n$ amoebots is placed on the triangular grid. An amoebot is able to send information to its neighbors, and to move via expansions and contractions. Since amoebots and information can only travel node by node, most problems have a natural lower bound of $\Omega(D)$ where $D$ denotes the diameter of the structure. Inspired by the nervous and muscular system, Feldmann et al. have proposed the reconfigurable circuit extension and the joint movement extension of the amoebot model with the goal of breaking this lower bound.

In the joint movement extension, the way amoebots move is altered. Amoebots become able to push and pull other amoebots. Feldmann et al. demonstrated the power of joint movements by transforming a line of amoebots into a rhombus within $O(\log n)$ rounds. However, they left the details of the extension open. The goal of this paper is therefore to formalize the joint movement extension. In order to provide a proof of concept for the extension, we consider two fundamental problems of modular robot systems: reconfiguration and locomotion.

We approach these problems by defining meta-modules of rhombical and hexagonal shapes, respectively. The meta-modules are capable of movement primitives like sliding, rotating, and tunneling. This allows us to simulate reconfiguration algorithms of various modular robot systems. Finally, we construct three amoebot structures capable of locomotion by rolling, crawling, and walking, respectively.


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## 1 Introduction

Programmable matter consists of homogeneous nano-robots that are able to change the properties of the matter in a programmable fashion, e.g., the shape, the color, or the density [57]. We are considering the geometric amoebot model $[14,15,16]$ where a set of $n$ nano-robots (called amoebots) is placed on the triangular grid. An amoebot is able to send information to its neighbors, and to move by first expanding into an unoccupied adjacent node, and then contracting into that node. Since amoebots and information can only travel node by node, most problems have a natural lower bound of $\Omega(D)$ where $D$ denotes the

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diameter of the structure. Inspired by the nervous and muscular system, Feldmann et al. [23] proposed the reconfigurable circuit extension and the joint movement extension with the goal of breaking this lower bound.

In the reconfigurable circuit extension, the amoebot structure is able to interconnect amoebots by circuits. Each amoebot can send a primitive signal on circuits it is connected to. The signal is received by all amoebots connected to the same circuit. Among others, Feldmann et al. [23] solved the leader election problem, compass alignment problem, and chirality agreement problem within $O(\log n)$ rounds w.h.p. ${ }^{1}$ These problems will be important to coordinate the joint movements. Afterward, Padalkin et al. [47] explored the structural power of the circuits by considering the spanning tree problem and symmetry detection problem. Both problems can be solved within polylogarithmic time w.h.p.

In the joint movement extension, the way amoebots move is altered. In a nutshell, an expanding amoebot is capable of pushing other amoebots away from it, and a contracting amoebot is capable of pulling other amoebots towards it. Feldmann et al. [23] demonstrated the power of joint movements by transforming a line of amoebots into a rhombus within $O(\log n)$ rounds. However, they left the details of the extension open. The goal of this paper is therefore to formalize the joint movement extension. In order to provide a proof of concept for the extension, we consider two fundamental problems of modular robot systems (MRS): reconfiguration and locomotion. We study these problems from a centralized view to explore the limits of the extension.

In the reconfiguration problem, an MRS has to reconfigure its structure into a given shape. Examples of reconfiguration algorithms in the original amoebot model can be found in $[17,18,35,40]$. However, all of these are subject of the aforementioned natural lower bound. To our knowledge, polylogarithmic time solutions were found for two types of MRSs: in the nubot model [62] and crystalline atom model [6]. We will show that in the joint movement extension, the amoebots are able to simulate the reconfiguration algorithm for the crystalline atom model, and with that break the lower bound.

In the locomotion problem, an MRS has to move along an even surface as fast as possible. We might also ask the MRS to transport an object along the way. In the original amoebot model, one would use the spanning tree primitive to move along the surface [14]. However, we only obtain a constant velocity with that. Furthermore, the original amoebot model does not allow us to transport any objects. In terrestrial environments, there are three basic types of locomotion: rolling, crawling, and walking [26, 33]. For each of these locomotion types, we will present an amoebot structure.

## 2 Preliminaries

In this section, we introduce the geometric amoebot model [15, 16] and formalize the joint movement extension.

### 2.1 Geometric Amoebot Model

In this section, we introduce the geometric amoebot model [15]. We slightly deviate from the original model to make it more suitable to our extension. A set of $n$ amoebots is placed on the infinite regular triangular grid graph $G_{\Delta}=(V, E)$ (see Figure 1a). An amoebot is an

[^0]

Figure 1 Domains. The figures show the domains we are working with. The tesselations will be explained in Section 3. Note that there are spaces between the meta-modules since a node cannot be occupied by more than one amoebot. However, the spaces do not contain any nodes.
anonymous, randomized finite state machine in the form of a line segment. The endpoints may either occupy the same node or two adjacent nodes. If the endpoints occupy the same node, the amoebot has length 0 and is called contracted and otherwise, it has length 1 and is called expanded. Every node of $G_{\Delta}$ is occupied by at most one amoebot. Two endpoints of different amoebots that occupy adjacent nodes in $G_{\Delta}$ are connected by bonds (red edges). Let the amoebot structure $S \subseteq V$ be the set of nodes occupied by the amoebots. We say that $S$ is connected iff $G_{S}$ is connected, where $G_{S}=\left.G_{\Delta}\right|_{S}$ is the graph induced by $S$. Initially, $S$ is connected. An amoebot can move through contractions and expansions. We refer to [15] for more details.

### 2.2 Joint Movement Extension

In the joint movement extension [23], the way the amoebots move is altered. The idea behind the extension is to allow amoebots to push and pull other amoebots. In the following, we formalize the joint movement extension.

We assume the fully synchronous activation model, i.e., the time is divided into synchronous rounds, and every amoebot is active in each round. Furthermore, we make the idealistic assumption that all movements start at the same time and are performed at the same speed. W.l.o.g., we assume that all movements happen within the time period $[0,1]$.

Joint movements are performed in two steps. In the first step, the amoebots remove bonds from $G_{S}$ as follows. Each amoebot can decide to release an arbitrary subset of its currently incident bonds in $G_{S}$. A bond is removed iff either of the amoebots at the endpoints releases the bond. However, an expanded amoebot cannot release the bond connecting its occupied nodes. Let $E_{L} \subseteq E_{S}$ be the set of all line segments (amoebots) of length $1, E_{R} \subseteq E_{S}$ be the set of the remaining bonds, and $G_{R}=\left(S, E_{L} \cup E_{R}\right)$ be the resulting graph. We require that $G_{R}$ is connected since otherwise, disconnected parts might float apart. We say that a connectivity conflict occurs iff $G_{R}$ is not connected. Whenever a connectivity conflict occurs, the amoebot structure transitions into an undefined state such that we become unable to make any statements about the structure.

In the second step, each amoebot may perform one of the following movements. A contracted amoebot may expand on one of the axes as follows (see blue amoebot in Figure 2a). At $t=0$, the amoebot can reorientate itself and reassign each of its incident bonds to one of its endpoints. Bonds assigned to an endpoint will stay with that endpoint as the amoebot expands. At $t \in[0,1]$, the amoebots has a length of $t$. In the process, the incident bonds do not change their orientations or lengths. As a result, the expanding amoebot pushes all connected amoebots. An expanded amoebot may contract analogously by reversing the contraction (see green amoebot in Figure 2b). Thereby, it pulls all connected amoebots.

(a) Expansion.

(b) Contraction.

(c) Handover.

Figure 2 Movements in the extension. Red lines indicate bonds. Blue amoebots are expanding. Green amoebots are contracting. The first two figures show a movement in 0.5 time steps.

(a) Structural conflicts and mappings. In the left figure, the expansions cause a structural conflict. Hence, the amoebot structure transitions into an undefined state. In the right figure, the amoebot structure can be mapped onto the triangular grid after the movement.

(b) Collision. Initially, we have a valid amoebot structure given $(t=0)$. At $t=1$, the amoebot structure could be mapped on $G_{\Delta}$. However, for $t \in[0.25,0.75]$, parts of the structure collide. Hence, the amoebot structure transitions into an undefined state.

Figure 3 Joint movements. Red lines indicate bonds. Blue amoebots are expanding horizontally.

Furthermore, a contracted amoebot $x$ occupying node $u$ and an expanded amoebot $y$ occupying nodes $v$ and $w$ may perform a handover if there is a bond $b$ between $u$ and $v$, as follows (see Figure 2c where $x$ is marked in blue and $y$ is marked in green). At an arbitrary $t \in[0,1]$, we switch the association of the endpoint $v$ from $y$ to $x$ such that $x$ becomes an expanded amoebot with endpoints occupying nodes $u$ and $v, y$ becomes a contracted amoebot with both endpoints occupying $w$, and $x$ and $y$ are connected by bond $\{v, w\}$. We have to include the handover to ensure universality of the model since otherwise, it would not be possible to move through a narrow tunnel. In theory, all movement primitives presented in Section 3 can be realized without handovers. However, for reasons of clarity and comprehensibility, we will still make use of the handovers.

The amoebots may not be able to perform their movements. We distinguish between two cases. First, the amoebots may not be able to perform their movements while maintaining their relative positions (see Figure 3a). We call that a structural conflict. Second, parts of the structure may collide into each other. More precisely, a collision occurs if there is a $t \in[0,1]$ such that two non-adjacent bonds intersect at some point (see Figure 3b). Whenever either a structural conflict or a collision occurs, the amoebot structure transitions into an undefined state such that we become unable to make any statements about the structure.

Otherwise, at $t=1$, we map the amoebots on the triangular grid $G_{\Delta}$ in compliance with the orientations of all bonds and line segments (see Figure 3a). The mapping is unique except for translations since $G_{R}$ is connected. We choose any of these mappings as our next amoebot structure. Let $S^{\prime}$ be the set of nodes occupied the amoebots, $E_{L}^{\prime}$ be the set of all line segments (amoebots) of length 1 after all movements were completed, and $G_{M}=\left(S^{\prime}, E_{L}^{\prime} \cup E_{R}\right)$ be the resulting graph. Afterwards, the amoebots reestablish all possible bonds, i.e., we obtain the graph $G_{S^{\prime}}=\left.G_{\Delta}\right|_{S^{\prime}}$ induced by $S^{\prime}$. Unless stated otherwise, each arrow in the figures indicates a single round. The left side shows the structure after the removal of bonds (first step), and the right side the structure after the execution of movements (second step). We chain multiple rounds if we do not change bonds.

In this paper, we assume that we have a centralized scheduler. The scheduler knows the current state of the amoebot structure at all times. At the beginning of each synchronous round, it decides for each amoebot (i) which bonds to release and (ii) which movements to perform. We leave the design of distributed solutions for future work.

### 2.3 Problem Statement and Our Contribution

In this paper, we formalize the joint movement extension proposed by Feldmann et al. [23]. In the following, we provide a proof of concept. For that, we focus on two fundamental problems of MRSs: reconfiguration and locomotion. We study these problems from a centralized view to explore the limits of the extension.

In the reconfiguration problem, an MRS has to reconfigure its structure into a given configuration. For that, we define meta-modules of rhombical and hexagonal shape. We show that these meta-modules are able to perform various movement primitives of other MRSs, e.g., crystalline atoms, and rectangular/hexagonal metamorphic robots. This allows us to simulate the reconfiguration algorithms of those models.

In the locomotion problem, an MRS has to move along an even surface as fast as possible. We might also ask the MRS to transport an object along the way. We present three amoebot structures that are able to move by rolling, crawling, and walking, respectively. We analyze their velocities and compare them to other structures of similar models.

### 2.4 Related Work

MRSs can be classified into various types, e.g., lattice-type, chain-type, and mobile-type $[1,9,66,67]$. We refer to the cited papers for examples. We will focus on lattice-type MRSs. These in turn can be characterized by three properties: (i) the lattice, (ii) the connectivity constraint, and (iii) the allowed movement primitives [2].

Various MRSs have been defined for different lattices, e.g., [12, 62] utilize the triangular grid, and $[3,22,51]$ utilize the Cartesian grid. We will build meta-modules for the Cartesian and the triangular grid. Note that some MRSs were also physically realized. Examples for MRSs using the triangular grid are hexagonal metamorphic robots [12], HexBots [52], fractal machines [42], and catoms [34]. Examples for MRSs using the Cartesian grid are CHOBIE II [56], EM-Cubes [7], M-Blocks [49], pneumetic cellular robots [28], and XBots [61].

We can identify four types of connectivity constraints: (i) the structure is connected at all times, (ii) the structure is connected except for moving robots, (iii) the structure is connected before and after movements, and (iv) there are no connectivity constraints. Our joint movement extension falls into the first category. Other examples are crystalline atoms [50], telecubes [55, 59], and prismatic cubes [60]. Examples for the second category are the sliding cube model [10, 24], rectangular [22] and hexagonal metamorphic robots [12], and for the third category the nubot model [62] and the line pushing model [3]. An example for the last category is the variant of the amoebot model considered by Dufoulon et al. [21].

The most common movement primitives for MRSs on the Cartesian grid are the rotation and slide primitives (see Figures 4 b and 4 c ), and for MRSs on the triangular grid the rotation primitive (see Figure 4a). Some models may constrain these movement primitives (see Figure 5). We refer to $[2,27]$ for a deeper discussion about possible constraints. One way to bypass such constraints is to construct meta-modules, e.g., see [19, 27, 43]. Our meta-modules implement all aforementioned movement primitives without any constraints, i.e., they perform them in place.


Figure 4 Examples of movement primitives. The green modules are moving, respectively. Models that utilize the Cartesian grid usually assume square modules. In Section 3.1, we utilize rhombical modules instead.

(a) Sliding model, e.g., [27].

(b) Pivoting model, e.g., [2].

Figure 5 Constrained movement primitives. The green modules are moving. The gray cells must be empty. Both primitives have the same result while the pivoting model requires more free space than the sliding model.

From all aforementioned models, crystalline atoms [50], telecubes [55, 59], and prismatic cubes [60] are the closest to our joint movement extension. In these MRSs, each robot has the shape of a unit square and is able to expand an arm from each side by half a unit. Adjacent robots are able to attach and detach their arms. Similar to our joint movement extension, a robot can move attached robots by expanding or contracting its arms. In contrast to amoebots, a line of robots cannot reconfigure to any other shape since each pair of robots can only have a single point of contact [5]. In order to allow arbitrary reconfigurations, the robots are combined into meta-modules of square shape that are capable of various movement primitives, e.g., the rotation, slide, contraction, and tunnel primitives (see Figures 4b-4e). Our rhombical meta-modules can implement all these movement primitives. The rhombical shape has two advantages compared to the square shape. First, we can implement further movement primitives that provide us with a simple way to construct a walking structure. Second, we can utilize rhombical meta-modules as a basis for hexagonal meta-modules.

## 3 Meta-Modules

In this section, we will combine multiple amoebots to meta-modules. In other models for programmable matter and modular robots, meta-modules have proven to be very useful. For example, they allow us to bypass restrictions on the reconfigurability [19,58] and to simulate (reconfiguration) algorithms for other models [4, 48]. In the subsections, we will present meta-modules of rhombical and hexagonal shape, respectively.

### 3.1 Rhombical Meta-Modules

Let $\ell$ be a positive even integer. Our rhombical meta-module consists of $\ell^{2} / 2$ uniformly oriented expanded amoebots that we arrange into a rhombus of side length $\ell-1$ (see Figure 6). We obtain a parallelogram of side lengths $\ell-1$ and $\ell / 2-1$ if we contract all amoebots (see Figure 6a). Note that we have to remove some bonds to perform the contraction. We can expand the parallelogram again by reversing the contractions.

(a) Contraction primitive.

(c) Realignment primitive.

(e) Slide primitive.

(f) $k$-tunnel primitive.

Figure 6 Movement primitives for rhombical meta-modules. Red meta-modules perform a pull operation, and blue meta-modules a push operation.

- Lemma 1. Our implementation of the contraction and expansion primitive requires a single round, respectively.

There are exactly two possibilities to arrange the uniformly oriented expanded amoebots in a rhombus. By reorienting the amoebots in pairs with the help of handovers, we can reorientate all amoebots within a rhombus (see Figure 6b).

- Lemma 2. Our implementation of the reorientation primitive requires 3 rounds.

Furthermore, there are three possibilities to align the sides of a rhombus to the axes of the triangular grid. By sliding each second row along its axis to the other side, we can realign the other axis a rhombus is aligned to (see Figure 6c). Note that in combination with the reorientation of the amoebots within a rhombus, we are able to align a rhombus with any two axes of the triangular grid.

- Lemma 3. Our implementation of the realignment primitive requires 2 rounds.

We can arrange the meta-modules on a rhombical tesselation of the plane if they are all aligned to the same axes (see Figure 1b). Note that due to the triangular grid, the meta-modules are not connected diagonally everywhere. Hence, we will only consider metamodules connected if their sides are connected. In the following, we introduce two movement primitives: the slide and $k$-tunnel primitive. Our implementations of these primitives are similar to the ones for crystalline robots (e.g., [5]) and teletubes (e.g., [59]).

In the slide primitive, we move a meta-module $R_{1}$ along two adjacent substrate metamodules $R_{2}$ and $R_{3}$ (see Figure 4c). We realize the primitive as follows (see Figure 6e). We assume that all amoebots are orientated into the movement direction. Otherwise, we apply the reorientation primitive. With respect to Figure 6e, let $L_{1}$ denote the uppermost layer of $R_{2}$ and $R_{3}$, and $L_{2}$ the second uppermost layer of $R_{2}$ and $R_{3}$. Our slide primitive consists
of two rounds. In the first round, we contract all amoebots in $L_{1}$ after removing all bonds between $L_{1}$ and $R_{1}$ except the last one in the movement direction, and all bonds between $L_{1}$ and $L_{2}$ except the last one in the opposite direction. This moves $R_{1}$ into its target position. In the second round, we restore $R_{2}$ and $R_{3}$. For that, we expand $L_{1}$ again after removing all bonds between $L_{1}$ and $R_{1}$, and between $L_{1}$ and $L_{2}$ except the last ones in the movement direction, respectively. This ensures that $R_{1}$ stays in place.

- Lemma 4. Our implementation of the slide primitive requires 2 rounds.

In the $k$-tunnel primitive, we move a meta-module $R$ through a simple path of metamodules with $k$ corners to the other end (see Figure 4e). However, we do not move $R$ directly through the path. Instead, we make use of the following two basic operations. First, by contracting two adjacent meta-modules into parallelograms, we pull one meta-module into the other (see the first round in Figure 6f). Second, by expanding two adjacent contracted meta-modules, we push one meta-module out again (see the fourth round in Figure 6f). Note that two adjacent contracted meta-modules form a rhombus consisting of $\ell^{2}$ contracted amoebots. Since the amoebots do not have any orientation, we can perform the push operation into any direction in parallel to the axes the meta-modules are aligned to.

Now, consider a line of at least 4 meta-modules with two contracted meta-modules at one end. By expanding those two meta-modules and contracting the two meta-modules at the other end, we transfer a meta-module from one end to the other end without changing the length of the line (see the third round in Figure 6f). If we have only a line of 3 meta-modules, it suffices to contract and expand one meta-module, respectively (see the second round in Figure 6 f ). Note that we have to remove most of the bonds along the line to permit the line to move freely. The pull and push operations allow us to transfer $R$ from one corner to the next corner in a single round.

- Lemma 5. Our implementation of the $k$-tunnel primitive requires $k+1$ rounds.

In particular, note that a 1-tunnel allows us to move a meta-module around another one (see Figure 6d). In other models for modular robot systems, this simple case is known as the rotation primitive.

- Lemma 6. Our implementation of the rotation primitive requires 2 rounds.


### 3.2 Hexagonal Meta-Modules

Let $\ell$ be an even integer as before. Our hexagonal meta-module consists of three rhombical meta-modules of side length $\ell-1$ (see Figure 7a). We arrange them into a hexagon of alternating side lengths $\ell$ and $\ell-1$.

There are two possibilities to arrange the rhombical meta-modules in the hexagon (see Figure 7a). We can switch between them as follows. Each rhombical meta-module within the hexagonal meta-module can be split into an equilateral triangle of side length $\ell-1$, and an equilateral triangle of side length $\ell-2$. The idea is to apply a similar technique as in the realignment primitive for rhombical meta-modules. In the first round, each rhombical meta-module contracts each second row without shifting the other rows. In the second round, the rhombical meta-modules interchange the smaller triangles through handovers. In the third round, each resulting rhombical meta-module expands each second row into the opposite direction (of the contracted amoebots of its bigger triangle) - again without shifting the other rows.

- Lemma 7. Our implementation of the switching primitive requires 3 rounds.


Figure 7 Movement primitives for hexagonal meta-modules.

We can arrange the meta-modules on a hexagonal tesselation of the plane (see Figure 1c). In the following, we introduce the rotation primitive for hexagonal meta-modules.

In the rotation primitive, we move a hexagonal meta-module $H_{1}$ around another hexagonal meta-module $H_{2}$ as follows (see Figure 7b). We arrange $H_{2}$ such that a rhombical metamodule $R_{2}$ is adjacent to both the old and new position of $H_{1}$, and $H_{1}$ such that the rhombical meta-module $R_{1}$ adjacent to $R_{2}$ is aligned to the same axes as $R_{2}$. We contract $R_{1}$ and $R_{2}$, and then expand them into the direction of the new position of $H_{1}$ (compare to Section 3.1). This movement primitive requires two rounds. Note that additional steps may be necessary to switch or reorientate the rhombical meta-modules beforehand.

- Lemma 8. Our implementation of the rotation primitive requires 2 rounds.


## 4 Reconfiguration

In this section, we discuss possible reconfiguration algorithms. For that, we look at reconfiguration algorithms for other lattice-type MRSs and discuss whether amoebots are capable of simulating these. In particular, we consider the polylogarithmic time solutions for the nubot model [62] and the crystalline atom model [6]. The first subsection deals with the former model while the second subsection deals with reconfiguration algorithms for our meta-modules including the one for the latter model.

### 4.1 Nubot Model

Similar to the amoebot model, the nubot model [62] considers robots on the triangular grid where at most one robot can be positioned on each node. Adjacent robots can be connected by rigid bonds ${ }^{2}$. In the joint movement extension, the rigid bonds correspond to $E_{R}$.

Robots are able to appear, disappear, and rotate around adjacent robots. A rotating robot may push and pull other robots into its movement direction. Hence, each rotation results in a translation of a set of connected robots (including the rotating robot) into the movement direction by the distance of 1 . The set depends on the bonds and the movement direction, and may include robots not connected by rigid bonds to the rotating robot. In

[^1]
(a) Amoebot $u \in C$ is adjacent to the same amoebot $v \in M$ before and after the translation.

(b) No amoebot in $C$ is adjacent to the same amoebot in $M$ before and after the translation.

Figure 8 Simulation of the nubot model. Let $d$ denote the northeastern direction, $d^{\prime}$ the southwestern direction, and $d^{\prime \prime}$ the northwestern direction. The blue amoebots indicate set $M$. The green amoebots indicate set $C$. The amoebots marked by a white dot indicate amoebots $u$ and $v$.
contrast to the joint movement extension, there are no collisions by definition of the set. However, a movement may not be performed due to structural conflicts. In this case, the nubot model does not perform the movement.

- Lemma 9. An amoebot structure of contracted amoebots is able to translate a set of robots in a constant number of joint movements if the translation is possible in the nubot model.

Proof. Let $d$ denote the movement direction, $d^{\prime}$ the opposite direction, and $d^{\prime \prime}$ any other cardinal direction. Let $M$ denote the set of amoebots that has to be moved. Instead of moving $M$ into direction $d$, we will move the remaining amoebot structure into direction $d^{\prime}$. Note that $M$ divides the remaining amoebot structure into connected components.
$\triangleright$ Claim 10. Each node $x$ in direction $d^{\prime}$ of an amoebot not in $M$ is either occupied by an amoebot of the same connected component or unoccupied

Proof. Trivially, $x$ cannot be occupied by an amoebot of another connected component. Further, $x$ cannot be occupied by an amoebot in $M$ since otherwise, the resulting amoebot structure would not be free of collisions.

For each connected component $C$, we perform the following steps in parallel. If there is an amoebot $u \in C$ that is adjacent to the same amoebot $v \in M$ before and after the translation, we proceed as follows (see Figure 8a). Let $A$ denote the row of amoebots in $C$ through $u$ into direction $d^{\prime \prime}$. Note that $A$ may only contain $u$. In the first move cycle, we remove all bonds between $C$ and $M$ except for the bond between $u$ and $v$, and each amoebot in $A$ expands into direction $d^{\prime}$. In the second move cycle, we remove all bonds between $C$ and $M$ except for the new bond between $u$ and $v$, and each amoebot in $A$ contracts. Note that both movements are possible due to Claim 10.

If there is no amoebot in $C$ that is adjacent to the same amoebot in $M$ before and after the translation, we proceed as follows (see Figure 8 b ). Let $B$ denote all amoebots in $C$ that have an unoccupied node in direction $d^{\prime}$. In the first move cycle, each amoebot in $B$ expands into direction $d^{\prime}$. There has to be an amoebot $u \in B$ that becomes adjacent to an amoebot $v \in M$. Otherwise, the resulting amoebot structure would not be connected. In the second move cycle, we remove all bonds between $C$ and $M$ except for the bond between $u$ and $v$, and each amoebot in $B$ contracts. Note that both movements are possible due to Claim 10.

Woods et al. [62] showed that in the nubot model, the robots are able to self-assemble arbitrary shapes/patterns in an amount of time equal to the worst-case running time for a Turing machine to compute a pixel in the shape/pattern plus an additional factor which is
polylogarithmic in its size. While we are able to perform the translations, we do not have the means to let amoebots appear and disappear in the amoebot model. This prevents us from simulating the reconfiguration algorithm by Woods et al. [62].

### 4.2 Reconfiguration Algorithms for Meta-Modules

Naturally, our meta-modules allow us to simulate reconfiguration algorithms for lattice-type MRSs of similar shape if we can implement the same movement primitives. This leads us to the following results.

- Theorem 11. There is a centralized reconfiguration algorithm for $m$ rhombical metamodules that requires $O(\log m)$ rounds and performs $\Theta(m \log m)$ moves overall.

Proof. Aloupis et al. [6] proposed a reconfiguration algorithm for crystalline atoms. It requires $O(\log m)$ rounds and performs $\Theta(m \log m)$ moves overall. The idea is to transform the initial shape to a canonical shape using a divide and conquer approach. The target shape is reached by reversing that procedure. We refer to [6] for the details. The algorithm utilizes the contraction, slide and tunnel primitives which our rhombical meta-modules are capable of (see Section 3.1). Hence, they can simulate this reconfiguration algorithm.

- Theorem 12. There is a centralized reconfiguration algorithm for $m$ hexagonal metamodules that requires $O(m)$ rounds. Each module has to perform at most $O(m)$ moves.

Proof. Hurtado et al. [27] proposed a reconfiguration algorithm for hexagonal robots. It requires $O(m)$ rounds and each module has to perform at most $O(m)$ moves. The idea is to compute a spanning tree and to move the robots along the boundary of the tree to a leader module where the robots form a canonical shape, e.g., a line. The target shape is reached by reversing that procedure. We refer to [27] for the details. The algorithm utilizes the rotation primitive which our hexagonal meta-modules are capable of (see Section 3.2). Hence, they can simulate this reconfiguration algorithm.

## 5 Locomotion

In this section, we consider amoebot structures capable of locomotion along an even surface. There are three basic types of terrestrial locomotion: rolling, crawling, and walking [26, 33]. We can find biological and artificial examples for each of those. In the following subsections, we will present an amoebot structure for each type and analyze their velocity. In the last subsection, we discuss the transportation of objects.

### 5.1 Rolling

Animals and robots that move by rolling either rotate their whole body or parts of it. Rolling is rather rare in nature. Among others, spiders, caterpillars, and shrimps are known to utilize rolling as a secondary form of locomotion during danger [8]. Bacterial flagella are an example for a creature that rotates a part of its body around an axle [39].

In contrast, rolling is commonly used in robotic systems mainly in the form of wheels. An example of a robot system rolling as a whole are chain-type MRSs that can roll by forming a loop, e.g., Polypod [63], Polybot [64, 65], CKBot [41, 53], M-TRAN [68, 38], and SMORES $[29,30]$. Further, examples for rolling robots can be found in $[8,9]$.

Our rolling amoebot structure imitates a continuous track that rotates around a set of wheels. Continuous tracks are deployed in various fields, e.g., construction, agriculture, and military. We build our continuous track structure from hexagonal meta-modules of

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Figure 9 Continuous track structure. The blue meta-modules rotate clockwise around the green meta-modules. We highlight one of the rotating meta-modules in a darker blue.
alternating side lengths $\ell$ and $\ell-1$ (see Figure 9). The structure consists of two parts: a connected substrate structure (green meta-modules), and a closed chain of meta-modules rotating along the outer boundary of the substrate (blue meta-modules).

The continuous track structure moves as follows. The rotating meta-modules that are in contact with the surface release all such bonds with the surface if they rotate away from the surface (see the dark blue meta-module in the third round). Otherwise, they keep these bonds such that the substrate structure is pushed forwards. Note that we have to apply the switching primitive between the rotations. We obtain the initial structure after two rotations. In doing so, the structure has moved a distance of $2 \cdot \ell$. By performing the movements periodically, we obtain the following theorem.

- Theorem 13. Our continuous track structure composed of hexagonal meta-modules of alternating side lengths $\ell$ and $\ell-1$ moves a distance of $2 \cdot \ell$ within each period of constant length.

Butler et al. [10] have proposed another rolling structure for the sliding-cube model that we are able to simulate. The structure resembles a swarm of caterpillars where caterpillars climb over each other from the back to the front [20]. However, due to stalling times, this structure is slower than our continuous track structure. We refer to [10] for the details.

(a) Simple line of amoebots.

(b) Line of rhombical meta-modules.

(c) Millipede structure. The red rhombi indicate the legs. The blue and green rhombi indicate the body. The blue rhombi indicate the connectors. Note that the figure only depicts the first half of a period.

Figure 10 Worm and millipede structures.

### 5.2 Crawling

Crawling locomotion is used by limbless animals. According to [31], crawling can be classified into three types: worm-like locomotion, caterpillar-like locomotion, and snake-like locomotion. We will explain the earthworm-like locomotion below and refer to [31] for the other two types. Due to the advantage of crawling in narrow spaces, various crawling structures have been developed for MRSs, e.g., crystalline atoms [36, 50, 51], catoms [11, 13], polypod [63], polybot [65, 69], M-TRAN [38, 68], and origami robots [32].

Our crawling amoebot structure imitates earthworms. An earthworm is divided into a series of segments. It can individually contract and expand each of its segments. Earthworms move by peristaltic crawling, i.e., they propagate alternating waves of contractions and expansions of their segments from the anterior to the posterior part. The friction between the contracted segments and the surface gives the worm grip. This anchors the worm as other segments expand or contract. The waves of contractions pull the posterior parts to the front, and the waves of expansions push the anterior parts to the front. [31, 44]

The simplest amoebot structure that imitates an earthworm is a simple line of $n$ expanded amoebots along the surface (see Figure 10a). Each amoebot can be seen as a segment of the worm. Instead of propagating waves of contractions and expansions, we contract and expand the whole structure at once. During each contraction (expansion), we release all bonds between the amoebot structure and the surface except for the ones at the head (tail) of the structure that serve as an anchor. As a result, the contraction (expansion) pulls (pushes) the structure to the front. The simple line has moved by a distance of $n$ along the surface after performing a contraction and an expansion. This is the fastest way possible to move along a surface since we accumulate the movements of all amoebots into the same direction. By performing the movements periodically, we obtain the following theorem.

- Theorem 14. A simple line of $n$ expanded amoebots moves a distance of $n$ every 2 rounds.

However, in practice, the contractions and expansions of the whole structure yield high forces acting on the connections within the amoebot structure. We can address this problem by thickening the worm structure. This increases the expansion of the structure and with

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that its stability. Consider a line of $r$ rhombical meta-modules of side length $\ell-1$ (see Figure 10b). Each module can be seen as a segment of the worm structure. Recall that we can contract a rhombical meta-module into a parallelogram (see Figure 6a). The line of rhombical meta-modules moves in the same manner as the simple line except for the following two points. First, we utilize the whole meta-module at the front and at the end as an anchor to increase the grip, respectively. Second, only the middle $r-2$ meta-modules participate in the contractions and expansions. The line of rhombical meta-modules moves a distance of $\frac{r-2}{2} \cdot \ell$ along the surface after performing a contraction and an expansion. By performing the movements periodically, we obtain the following theorem.

- Theorem 15. A line of r rhombical meta-modules of side length $\ell-1$ moves a distance of $\frac{r-2}{2} \cdot \ell$ every 2 rounds.

Another problem in practice is friction between the structure and the surface and with that the wear of the structure. The worm structure is therefore poorly scalable in its length such that other types for locomotion are more suitable for large amounts of amoebots.

Most of the cited MRSs at the beginning of this section are very similar to our construction. The construction for crystalline robots is the closest one. Each of these consists of a line of (meta-)modules that are able to contract. Katoy et al. [36] also propose a "walking" structure (see Figure 12). However, the locomotion is still caused by the contraction of the body instead of motions of the legs. So, it is rather a caterpillar-like crawling than a walking movement.

### 5.3 Walking

A wide variety of animals are capable of walking locomotion, e.g., mammals, reptiles, birds, insects, millipedes, and spiders. Just as wide is the variety of differences, e.g., they differ in the number of legs, in the structure of the legs, and in their gait. Walking structures have been built for chain-type MRSs, e.g., M-TRAN [68, 38] and polybot [65, 69].

Our walking amoebot structure imitates millipedes. Millipedes have flexible, segmented bodies with tens to hundreds of legs that provide morphological robustness [54]. They move by propagating leg-density waves from the posterior to the anterior [25, 37]. We build our millipede structure from rhombical meta-modules of side length $\ell-1$ (see Figure 10c). Let $p$ denote the number of legs. The body and each leg consists of a line of rhombical meta-modules. All legs have the same size. Let $q$ denote the number of rhombical modules in each leg. We attach each leg to a meta-module of the body and orientate them alternating to the front and to the back. We call those meta-modules connectors. In order to prevent the legs from colliding, we place $q$ meta-modules between two connectors. Altogether, the millipede structure consists of $(2 p-1) \cdot q+p$ meta-modules.

In order to move the legs back and forth, we simply apply the realignment primitive (see Figure 6c) on all meta-modules within the legs (see Figure 10c). We achieve forward motion by releasing all bonds between the surface and the legs moving forwards. In one step, the body moves a distance of $q \cdot \ell$. Note that the number of legs has no impact on the velocity. By continuously repeating these leg movements, we achieve a motion similar to the leg-density waves of millipedes. Note that we reach the initial amoebot structure after two leg movements. Hence, we obtain the following theorem.

- Theorem 16. Our millipede structure composed of rhombical meta-modules of side length $\ell-1$ with $p$ legs composed of $q$ rhombical meta-modules moves a distance of $2 \cdot q \cdot \ell$ within each period of constant length.

In practice, we can reduce the friction by additionally lifting the legs moving forwards (see Figure 13). For that, it suffices to partially contract all meta-modules of the body except for the connectors connected to a leg moving backwards. After the movement, we lower the lifted legs back to the ground. For that, we reverse the contractions within the body.

### 5.4 Transportation

Another important aspect of locomotion is the transportation of objects. The continuous track and worm structure are unsuitable for the transportation of objects due to their unstable top. In the worm structure, we can circumvent that problem by increasing the number of meta-modules not participating in the contractions and expansions. However, this decreases the velocity of the worm structure and increases the friction under the object.

In contrast, the millipede structure provides a rigid transport surface for the transportation of objects (see Figure 11). In practice, a high load introduced by the object on the structure may lead to problems. However, it was shown that millipedes have a large payload-to-weight ratio since they have to withstand high loads while burrowing in leaf litter, dead wood, or soil [25,54]. The millipede is able to distribute the weight evenly on its legs. Consequently, the millipede structure is well suited for the transportation of objects.


Figure 11 Transportation of objects by the millipede structure. The green rhombi indicate the millipede structure. The gray parallelogram indicates the object.

## 6 Conclusion and Future Work

In this paper, we have formalized the joint movement extension that were proposed by Feldmann et al. [23]. We have constructed meta-modules of rhombical and hexagonal shape that are able to perform various movement primitives. This allows us to simulate reconfiguration algorithms of various MRSs. However, our meta-modules are more flexible, e.g., we can move a hexagonal meta-module through two others. Such new movement primitives may lead to faster reconfiguration algorithms, e.g., sublinear solutions for hexagonal meta-modules or even arbitrary amoebot structures. Furthermore, we have presented three amoebot structures capable of moving along an even surface. In future work, movement on uneven surfaces can be considered.

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## A Omitted Figures

This section contains omitted figures.


Figure 12 Caterpillar structure. This structure is a replication of the structure by Katoy et al. [36] in the joint movement extension.


Figure 13 Movement of a millipede. For the sake of simplicity, we only show two legs.


[^0]:    ${ }^{1}$ An event holds with high probability (w.h.p.) if it holds with probability at least $1-1 / n^{c}$, where the constant $c$ can be made arbitrarily large.

[^1]:    ${ }^{2}$ Woods et al. [62] distinguish between rigid and flexible bonds. We ignore that since the flexible bonds are not necessary to achieve the polylogarithmic time reconfiguration algorithm.

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